

BLG 311E

Formal Languages and Automata

Recitation-2

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Deterministic and Non-Deterministic Finite Automata

Question

Design a DFA over the alphabet (a, b) that accepts strings of the form:

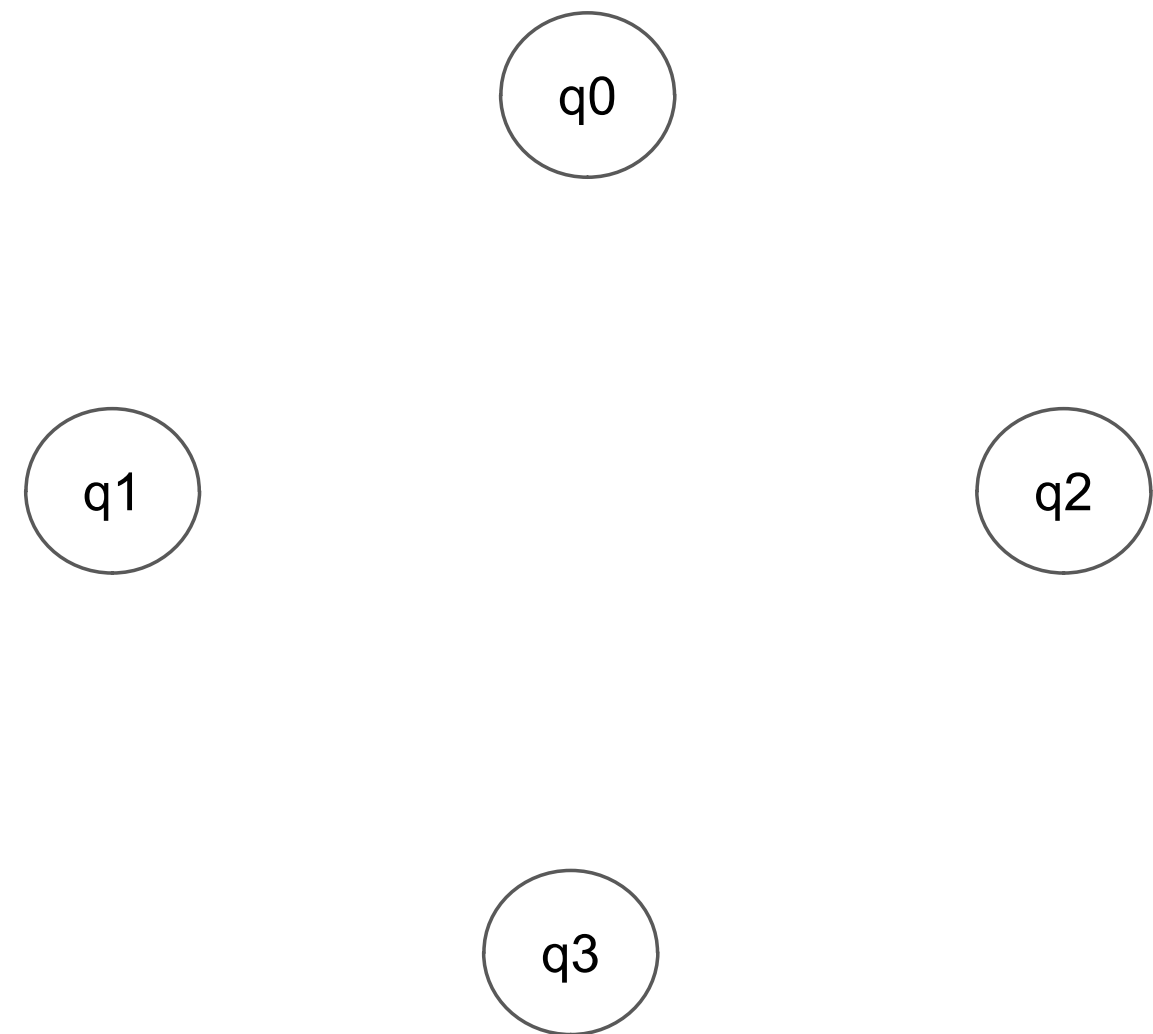
$$ab^{2k+1}, \quad k \geq 0$$

Solution

$$ab^{2k+1}, k \geq 0$$

States: q_0, q_1, q_2, q_3

Alphabets: a, b

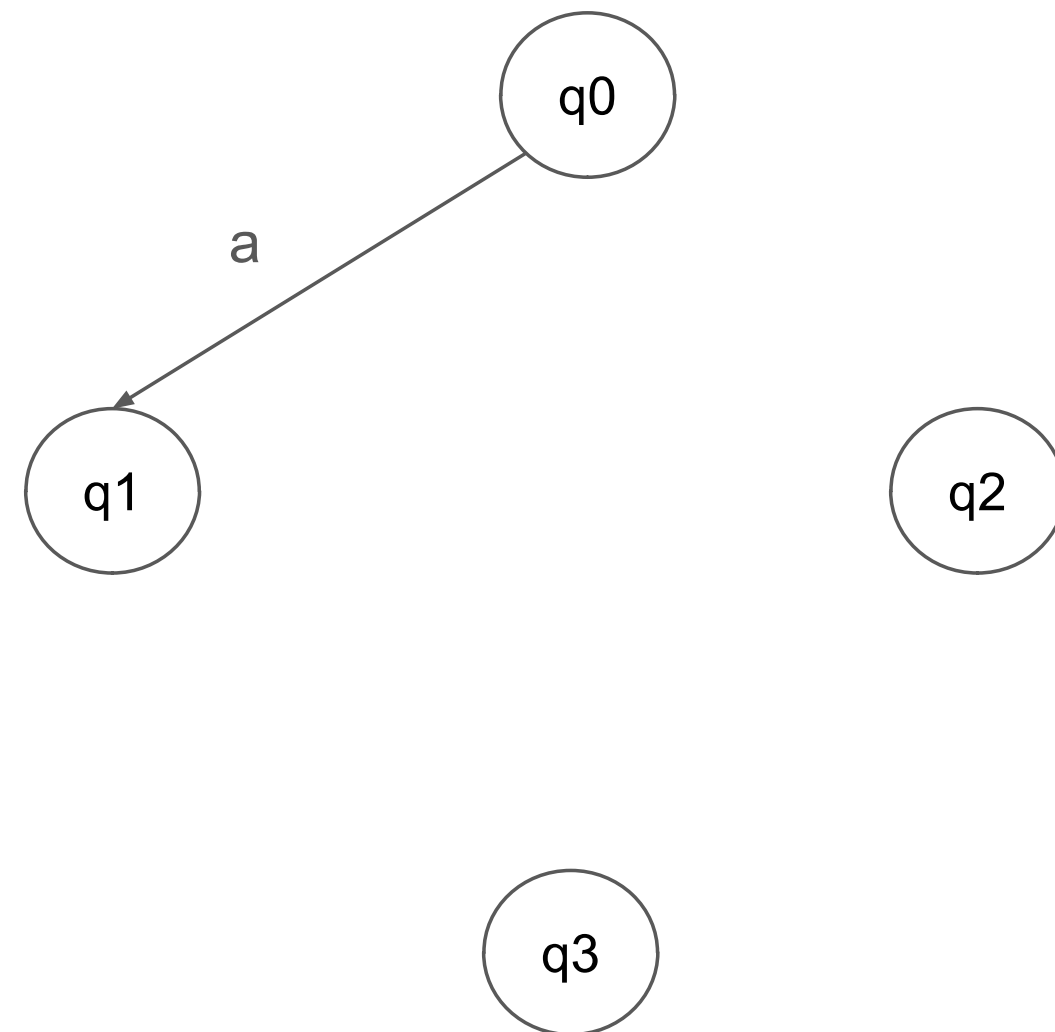


Solution

$ab^{2k+1}, k \geq 0$

States: q_0, q_1, q_2, q_3

Alphabets: $\{a, b\}$

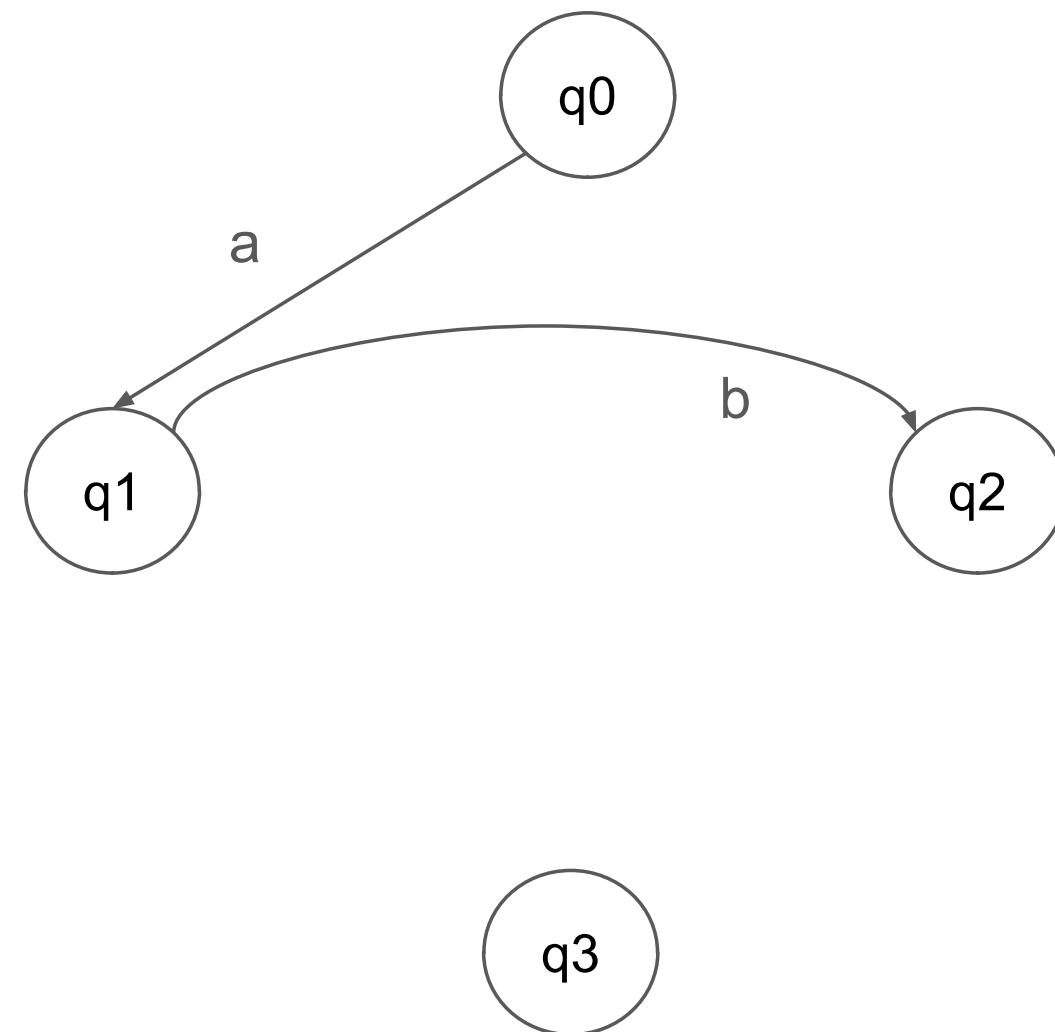


Solution

$$ab^{2k+1}, k \geq 0$$

States: Q0, Q1, Q2, Q3

Alphabets: a, b

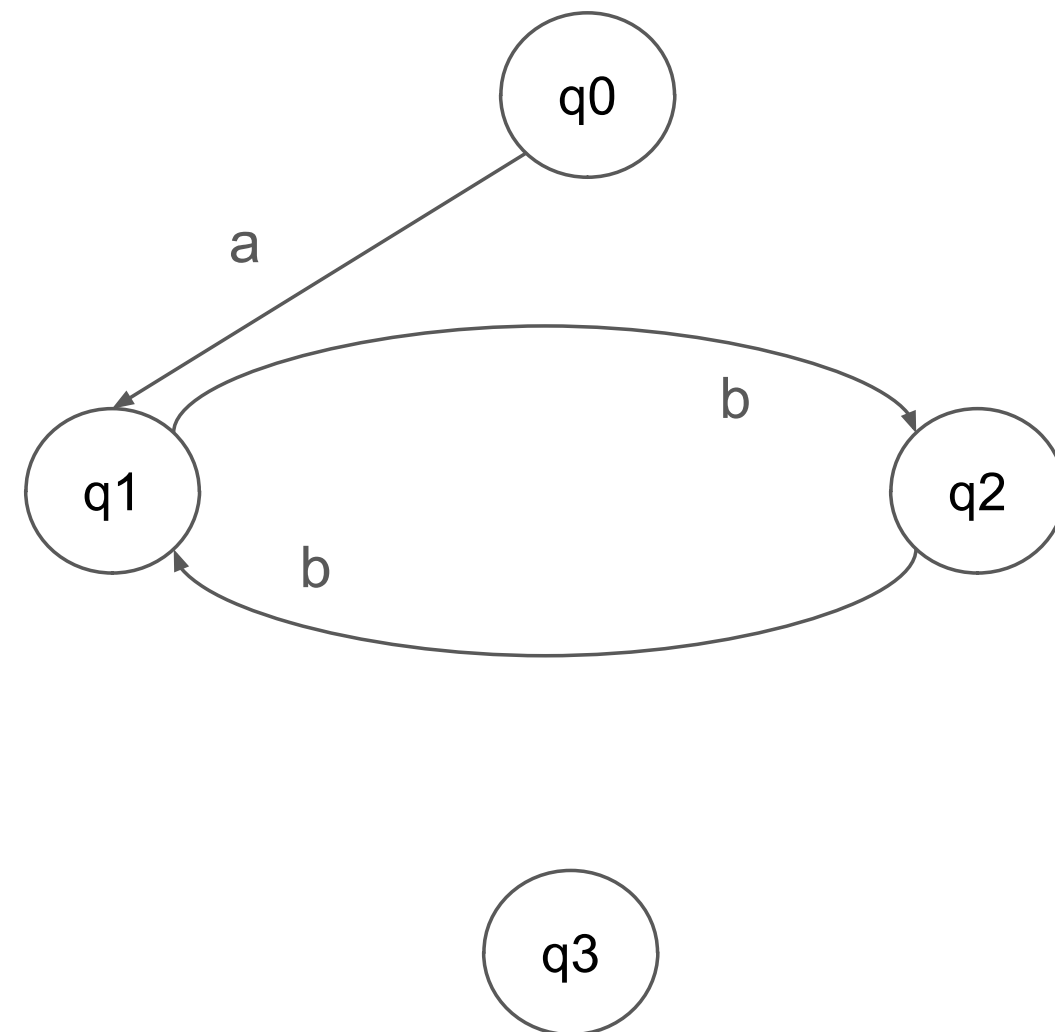


Solution

$$ab^{2k+1}, k \geq 0$$

States: Q0, Q1, Q2, Q3

Alphabets: a, b

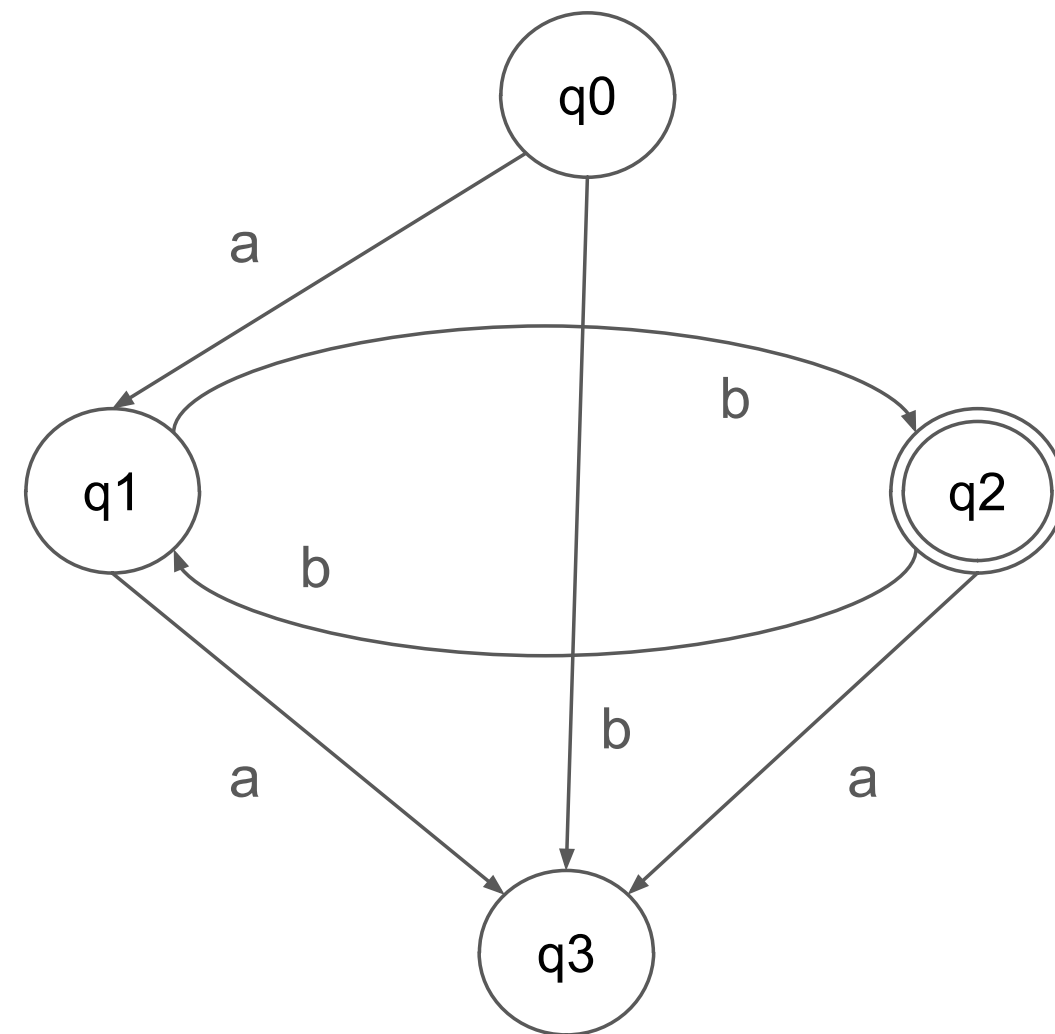


Solution

ab^{2k+1} , $k \geq 0$

States: Q0, Q1, Q2, Q3

Alphabets: a, b



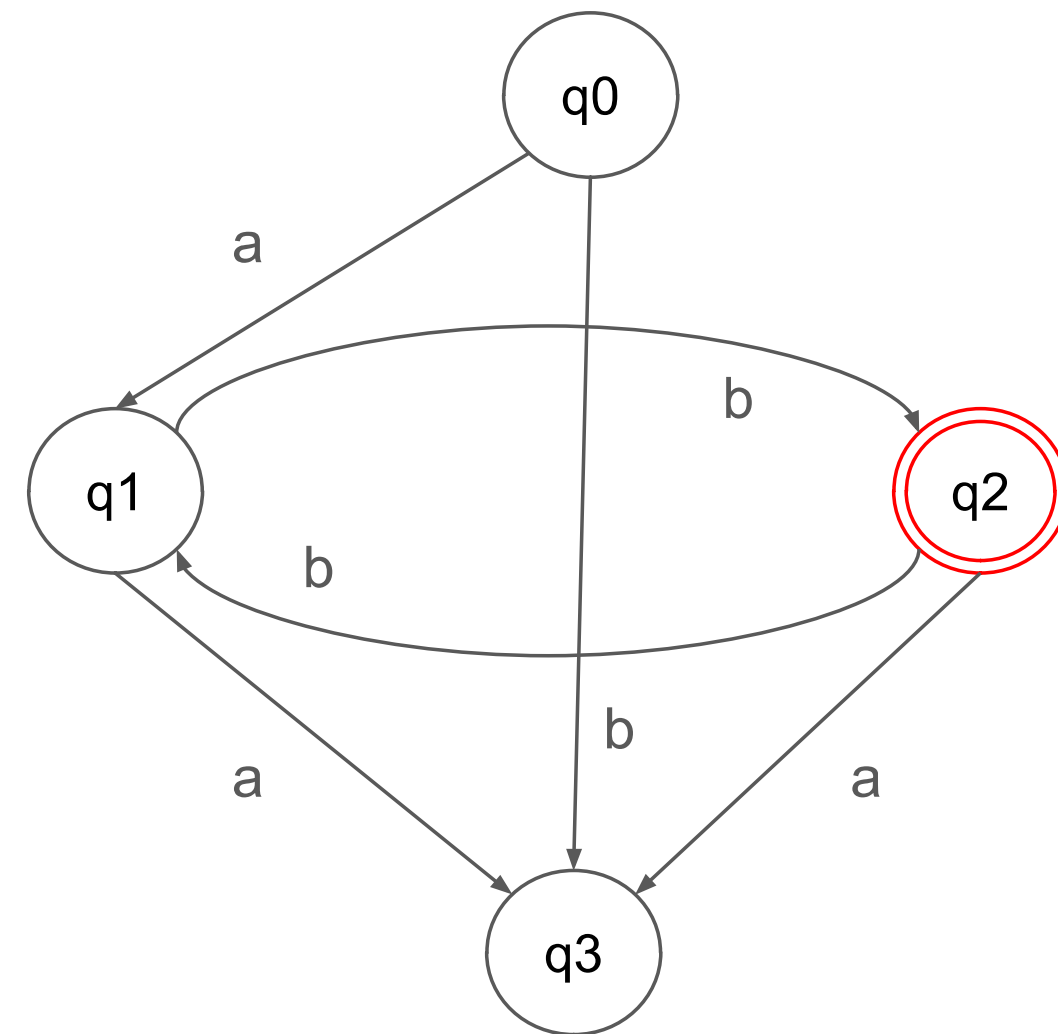
Solution

$$ab^{2k+1}, k \geq 0$$

States: Q0, Q1, Q2, Q3

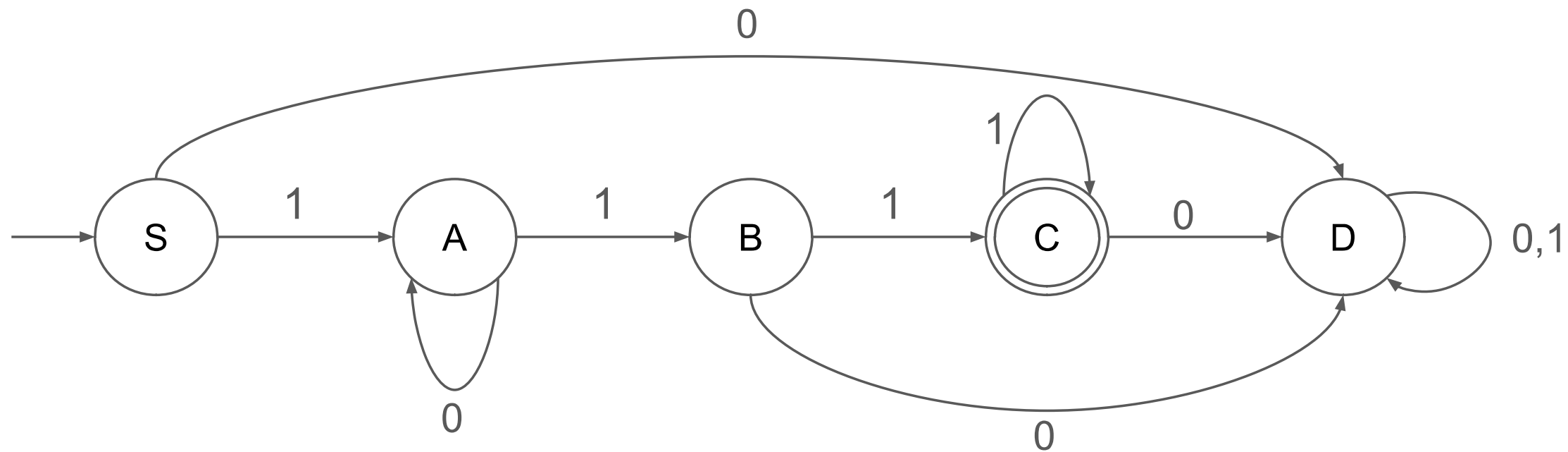
Alphabets: a, b

$ab(bb)^*$

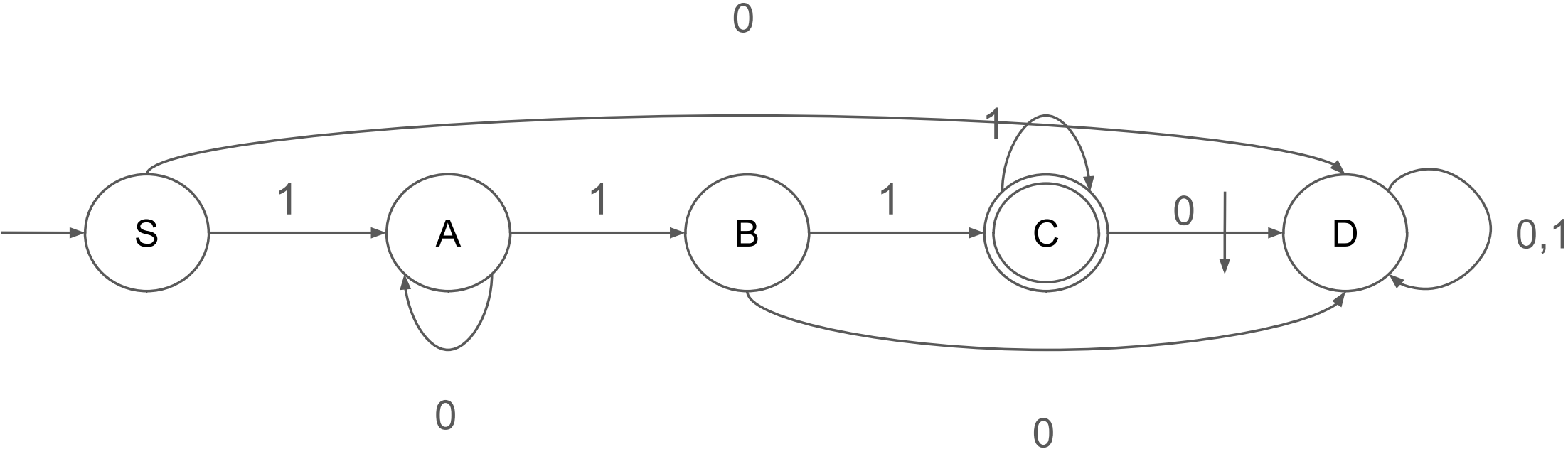


Question

Find the Grammar and regular expression of the given DFA automata below.

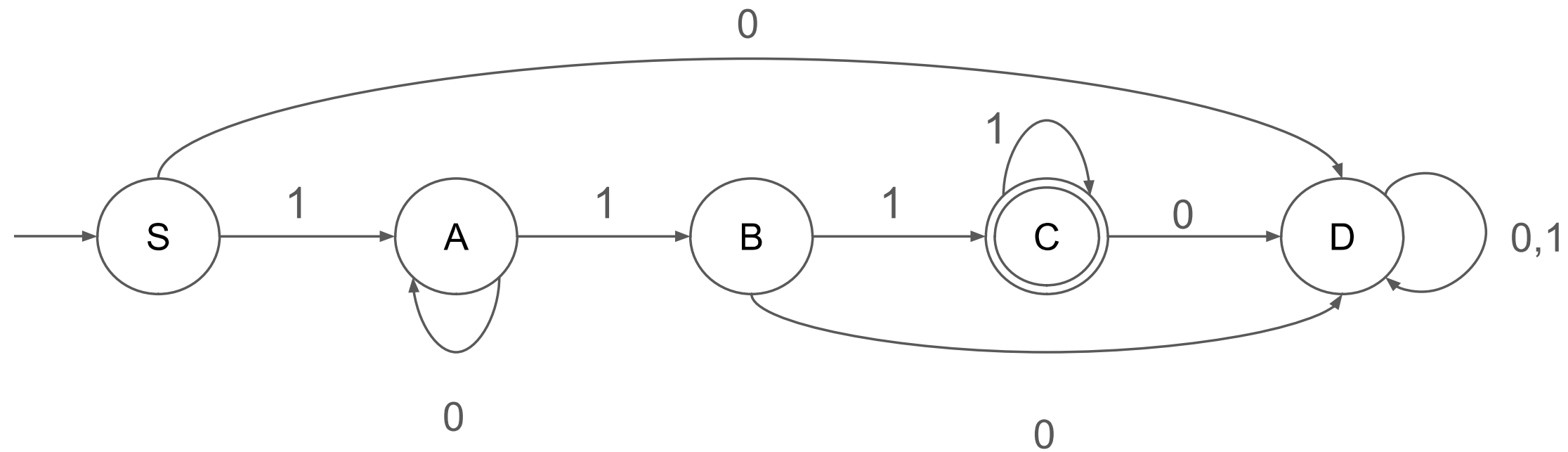


Solution



S→1A
A→0A|1B
B→1C
C→1C

Solution



$S \rightarrow 1A$

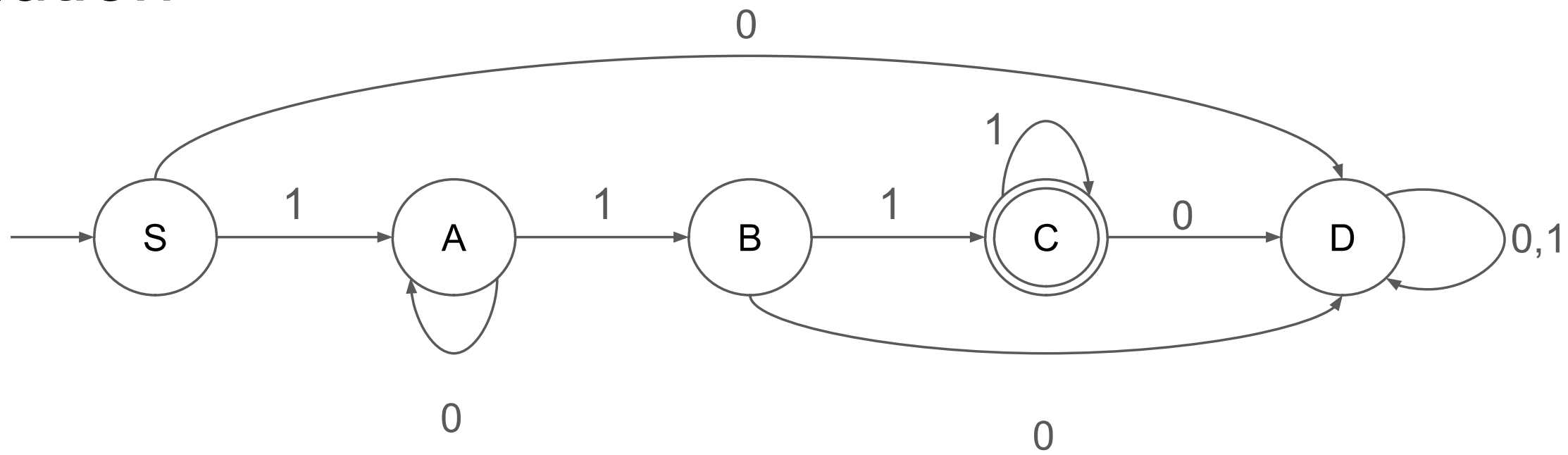
$A \rightarrow 0A \mid 1B$

$B \rightarrow 1C$

$C \rightarrow 1C$



Solution



$S \rightarrow 1A$

$A \rightarrow 0A \mid 1B$

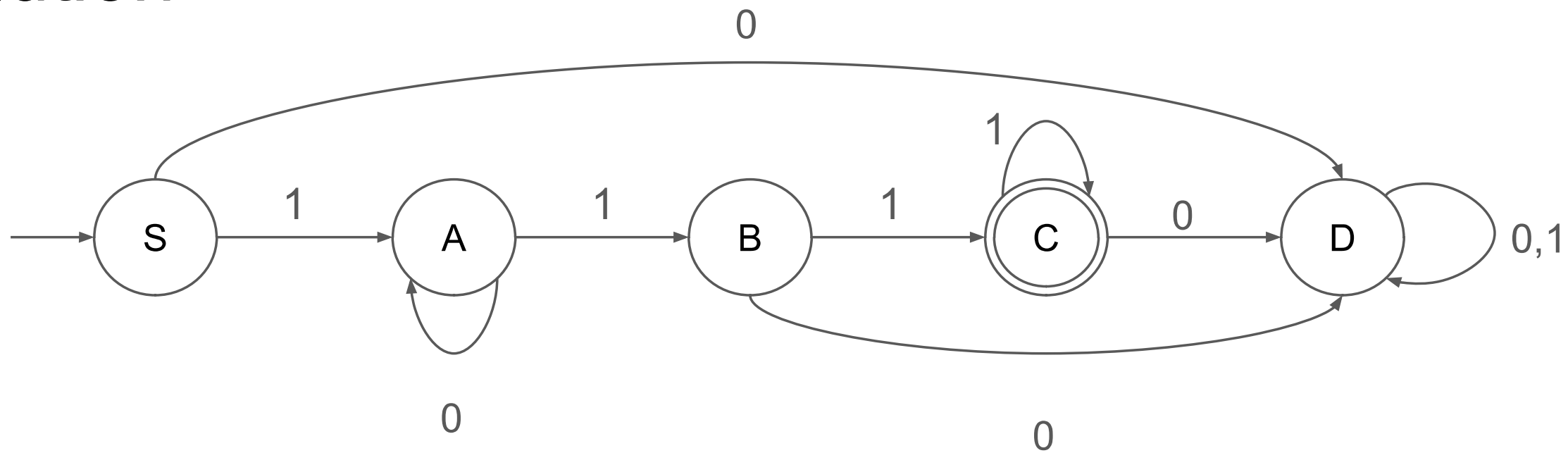
$B \rightarrow 1C$

$C \rightarrow 1C$



Regular Expression?

Solution



$S \rightarrow 1A$

$A \rightarrow 0A \mid 1B$

$B \rightarrow 1C$

$C \rightarrow 1C$



Regular Expression:

10^*111^*

Question

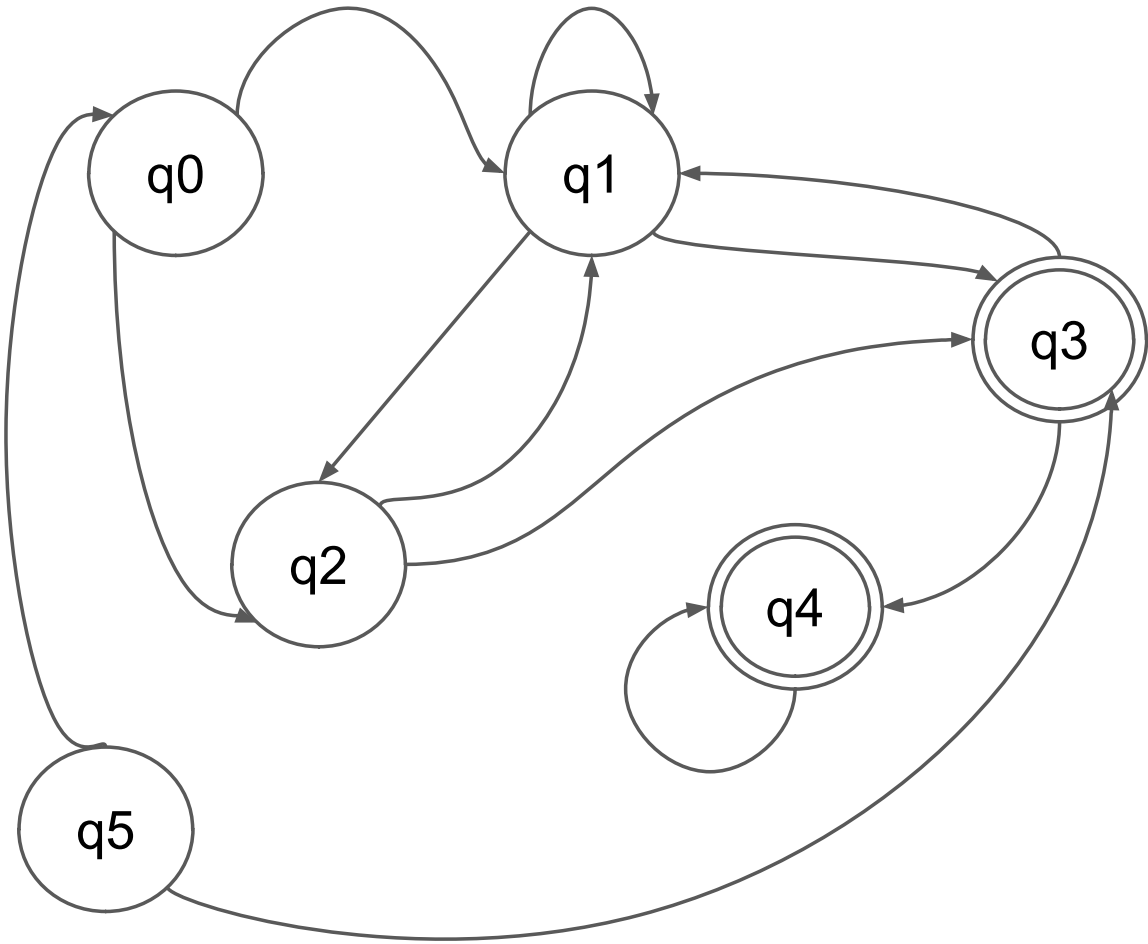
For the given transition table:

- A. Draw the Deterministic Finite Automaton (DFA) based on the transitions and output.
- B. Simplify the dependency table. Show all steps clearly.

State	Input a	Input b	Output
Q0	Q1	Q2	0
Q1	Q3	Q2	0
Q2	Q1	Q3	0
Q3	Q4	Q1	1
Q4	Q4	Q4	1
Q5	Q0	Q3	0

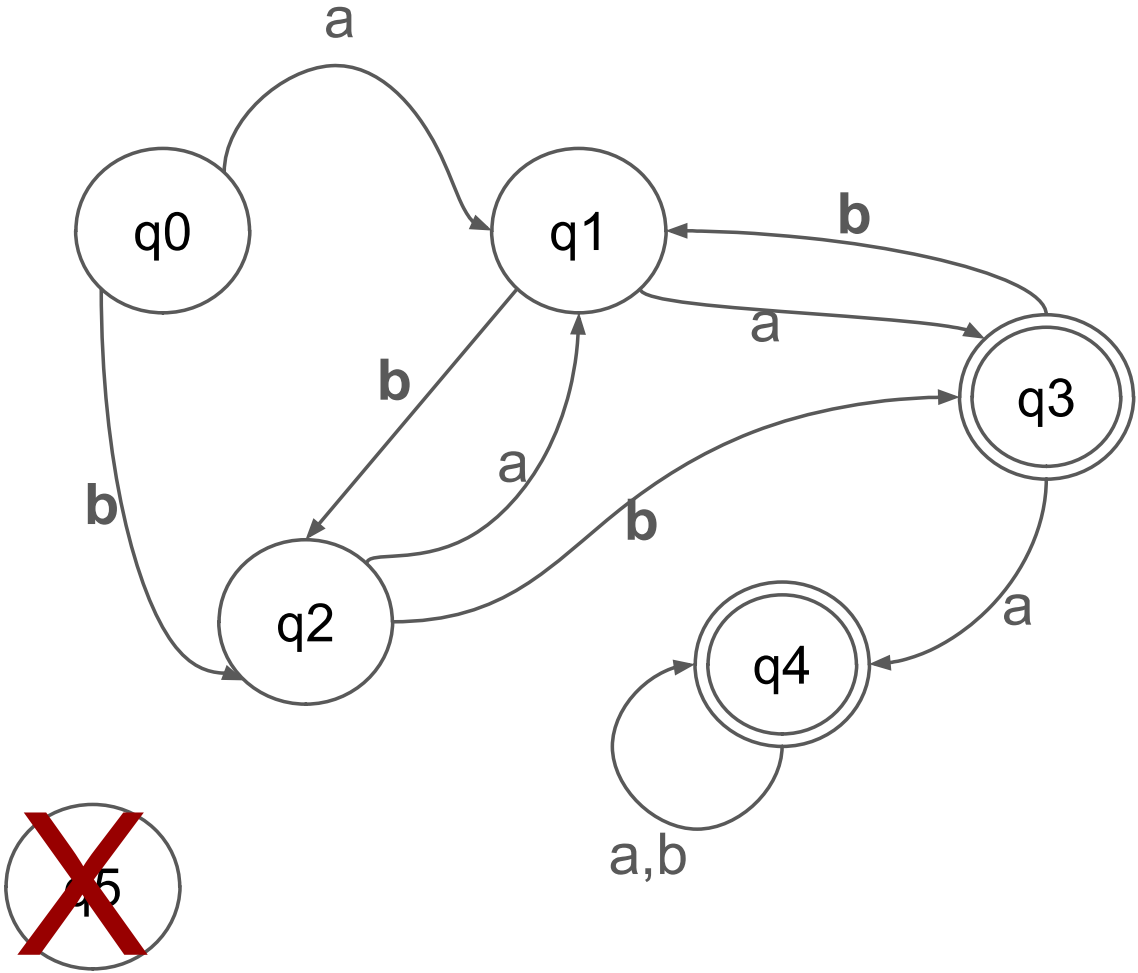
Solution

State	Input a	Input b	Output
Q0	Q1	Q2	0
Q1	Q3	Q2	0
Q2	Q1	Q3	0
Q3	Q4	Q1	1
Q4	Q4	Q4	1
Q5	Q0	Q3	0



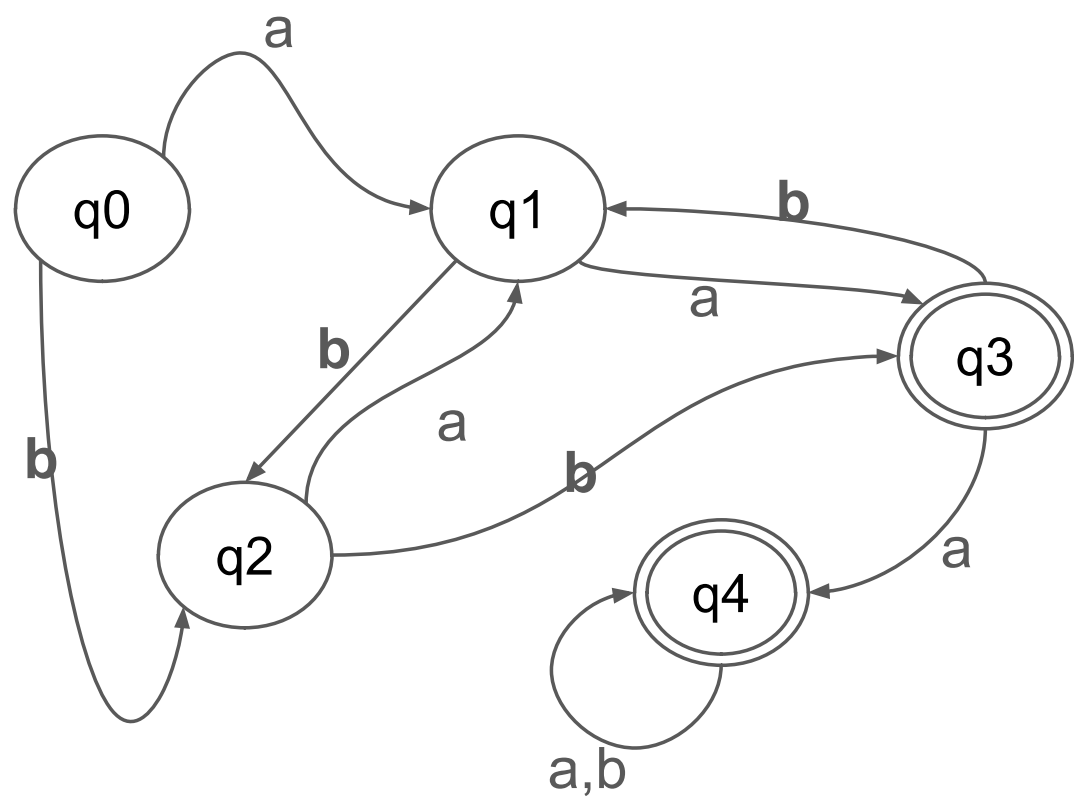
Solution

State	Input a	Input b	Output
Q0	Q1	Q2	0
Q1	Q3	Q2	0
Q2	Q1	Q3	0
Q3	Q4	Q1	1
Q4	Q4	Q4	1
Q5	Q0	Q3	0



Solution

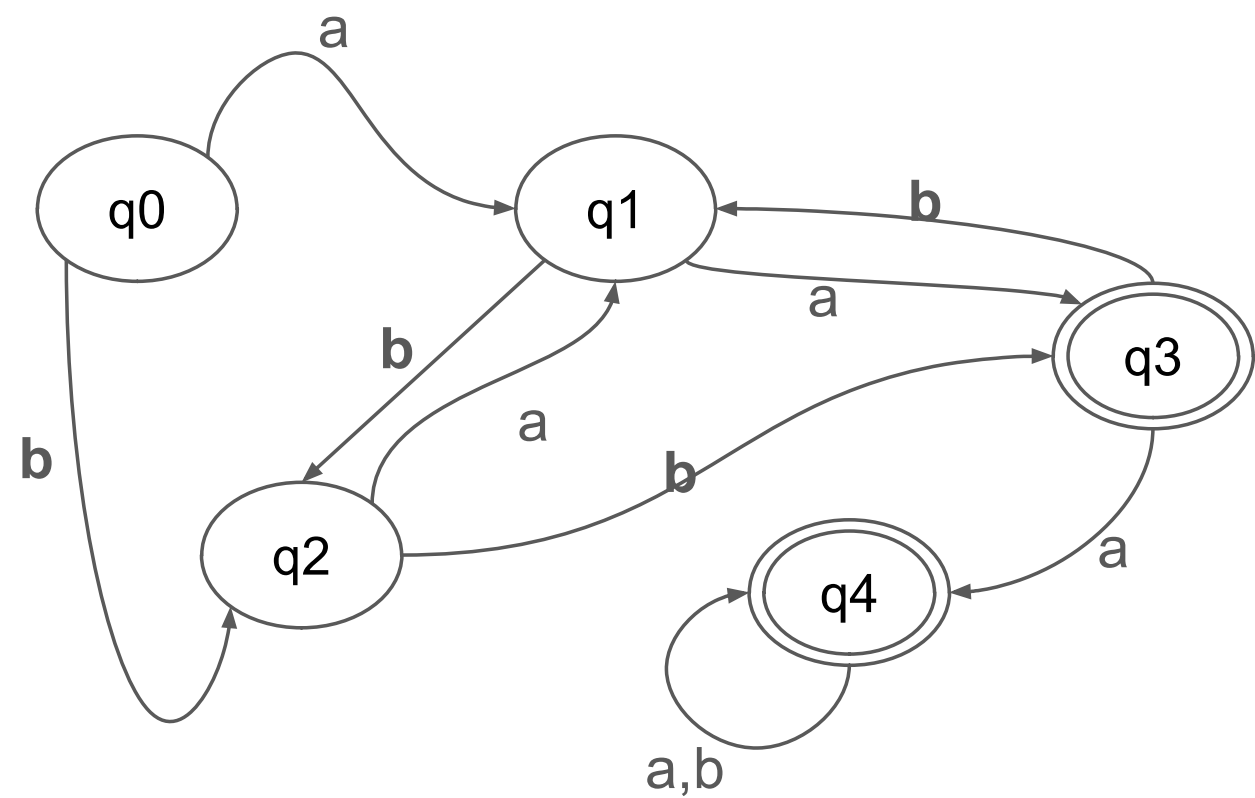
State	Input a	Input b	Output
Q0	Q1	Q2	0
Q1	Q3	Q2	0
Q2	Q1	Q3	0
Q3	Q4	Q1	1
Q4	Q4	Q4	1
Q5	Q0	Q3	0



	Q0	Q1	Q2	Q3	Q4
Q0	—				
Q1		—			
Q2			—		
Q3				—	
Q4					—

Solution

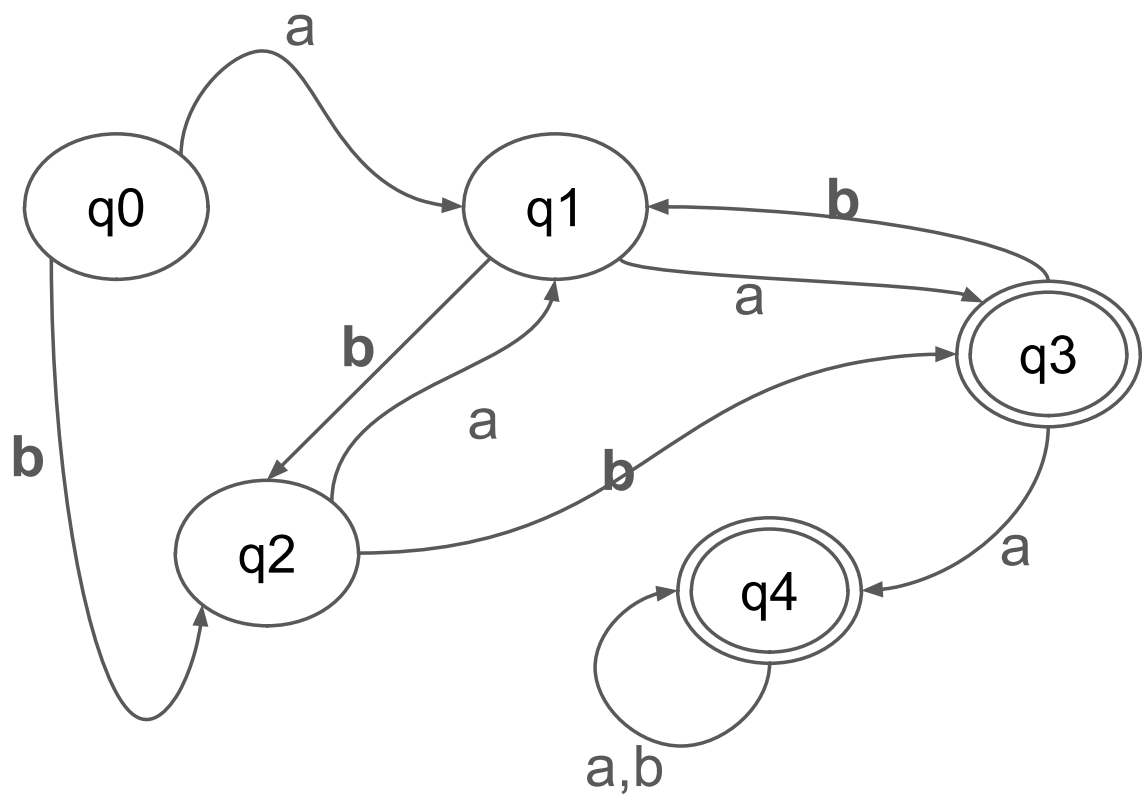
State	Input a	Input b	Output
Q0	Q1	Q2	0
Q1	<u>Q3</u>	Q2	0
Q2	Q1	<u>Q3</u>	0
<u>Q3</u>	<u>Q4</u>	Q1	1
<u>Q4</u>	<u>Q4</u>	<u>Q4</u>	1
Q5	Q0	Q3	0



	Q0	Q1	Q2	Q3	Q4
Q0	—	X	X	X	X
Q1		—	X	X	X
Q2			—	X	X
Q3				—	X
Q4					—

Solution

State	a	b	Output
Q0	Q1	Q2	0
Q1	Q3	Q2	0
Q2	Q1	Q3	0
Q3	Q4	Q1	1
Q4	Q4	Q4	1



	Q0	Q1	Q2	Q3	Q4
Q0	—	X	X	X	X
Q1		—	X	X	X
Q2			—	X	X
Q3				—	X
Q4					—

Question

Design an NFA over the alphabet $\{a, b\}$ that accepts all strings which:

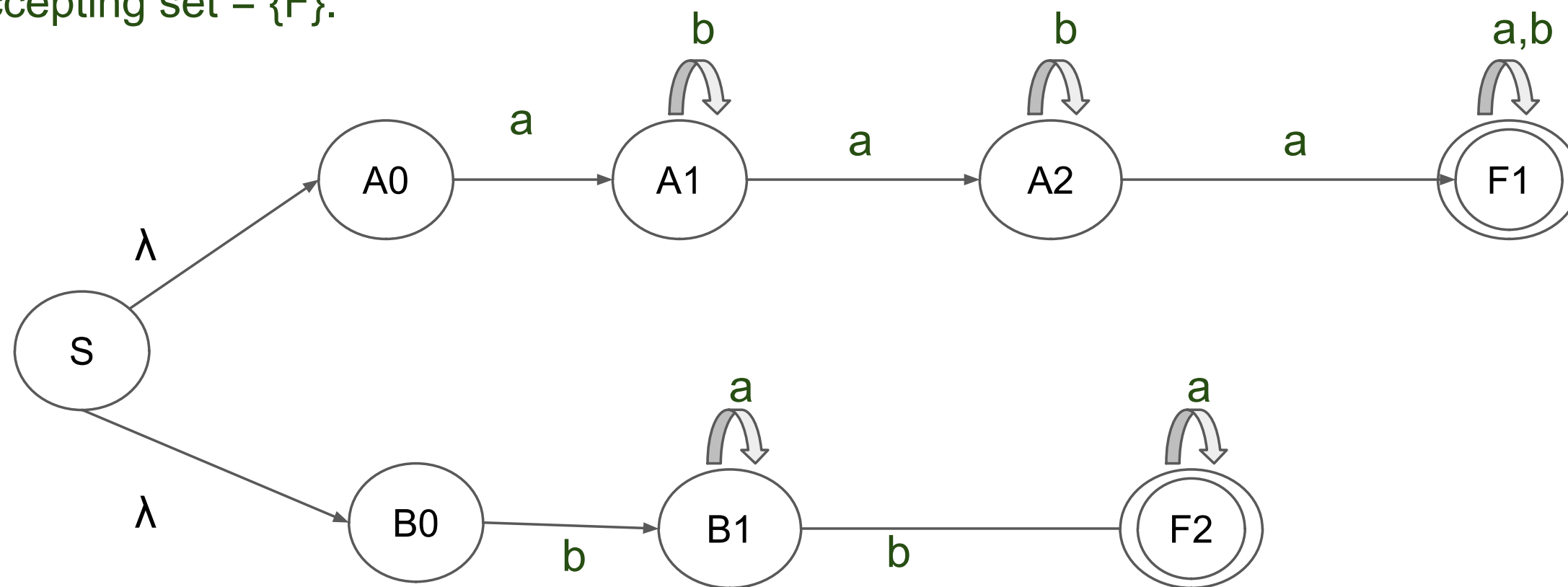
- Contain at least three a,
or
- Contain exactly two b.

Solution

Contain at least three “a” or Contain exactly two “b”.

Alphabet $\Sigma = \{a, b\}$.

Accepting set = $\{F\}$.



Question

Consider the following right-linear grammar over the alphabet $\{a, b\}$:

$$S \rightarrow aA \mid bB \mid a$$
$$A \rightarrow aS \mid b$$
$$B \rightarrow aA \mid bS$$

- A. Design a Non-deterministic Finite Automaton (NFA) that accepts the same language.
- B. Derive the regular expression corresponding to this language. Show your reasoning.

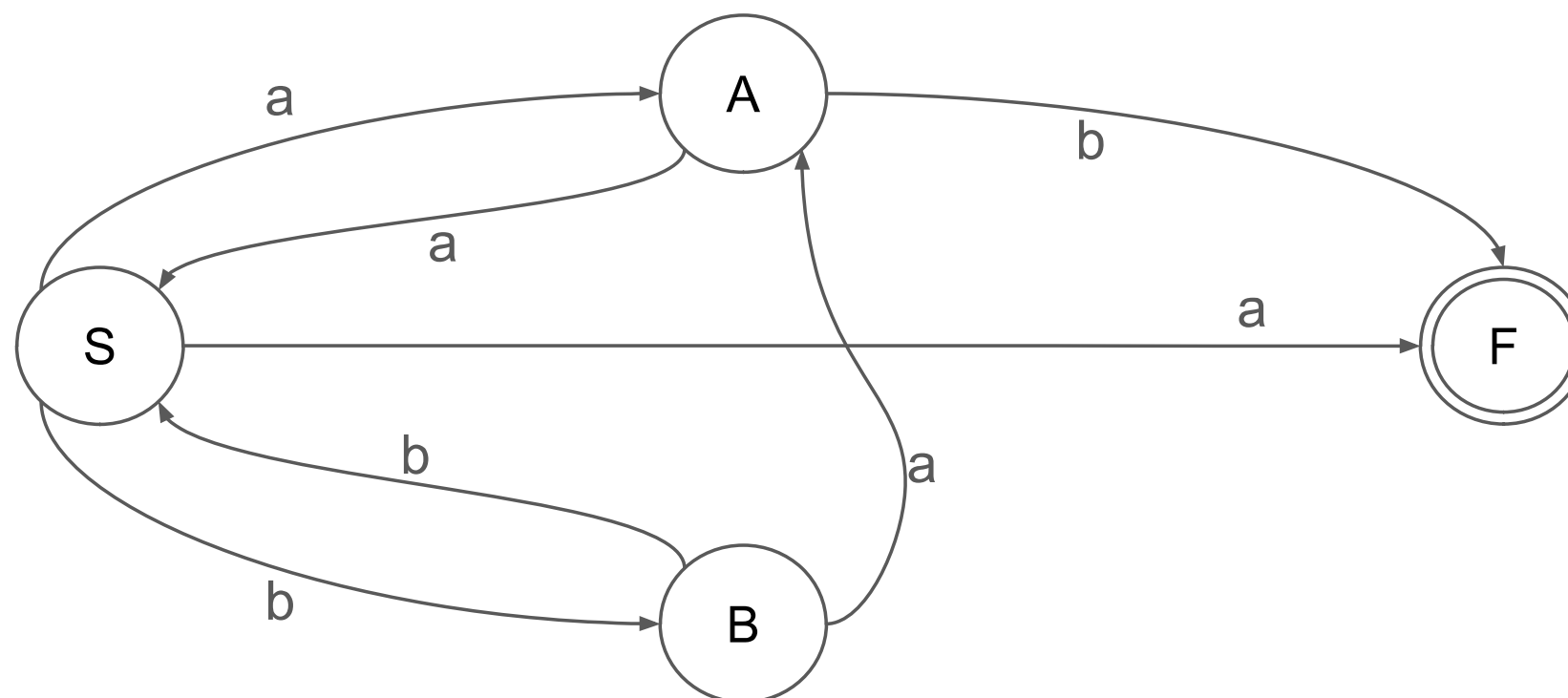
Solution

$S \rightarrow aA \mid bB \mid a$

$A \rightarrow aS \mid b$

$B \rightarrow aA \mid bS$

A. Design a Non-deterministic Finite Automaton (NFA) that accepts the same language.



Solution

$$S \rightarrow aA \mid bB \mid a$$

$$A \rightarrow aS \mid b$$

$$B \rightarrow aA \mid bS$$

B. Derive the regular expression corresponding to this language. Show your reasoning.

$$S = aA \cup bB \cup a$$

$$A = aS \cup b$$

$$B = aA \cup bS$$

Solution

$$S \rightarrow aA \mid bB \mid a$$

$$A \rightarrow aS \mid b$$

$$B \rightarrow aA \mid bS$$

B. Derive the regular expression corresponding to this language. Show your reasoning.

Eliminating A and B

$$S = aA \cup bB \cup a$$

$$A = aS \cup b$$

$$B = aA \cup bS$$

$$\begin{aligned} A = aS \cup b &\Rightarrow B = a(aS \cup b) \cup bS = aaS \cup ab \cup bS \\ \Rightarrow S &= a(aS \cup b) \cup b(aaS \cup ab \cup bS) \cup a \\ &= (aa \cup baa \cup bb)S \cup (a \cup ab \cup bab) \end{aligned}$$

Solution

$$S \rightarrow aA \mid bB \mid a$$

$$A \rightarrow aS \mid b$$

$$B \rightarrow aA \mid bS$$

B. Derive the regular expression corresponding to this language. Show your reasoning.

Eliminating A and B

$$S = aA \cup bB \cup a$$

$$A = aS \cup b$$

$$B = aA \cup bS$$

$$\begin{aligned} A = aS \cup b &\Rightarrow B = a(aS \cup b) \cup bS = aaS \cup ab \cup bS \\ \Rightarrow S &= a(aS \cup b) \cup b(aaS \cup ab \cup bS) \cup a \\ &= (aa \cup baa \cup bb)S \cup (a \cup ab \cup bab) \end{aligned}$$

Apply Arden's rule

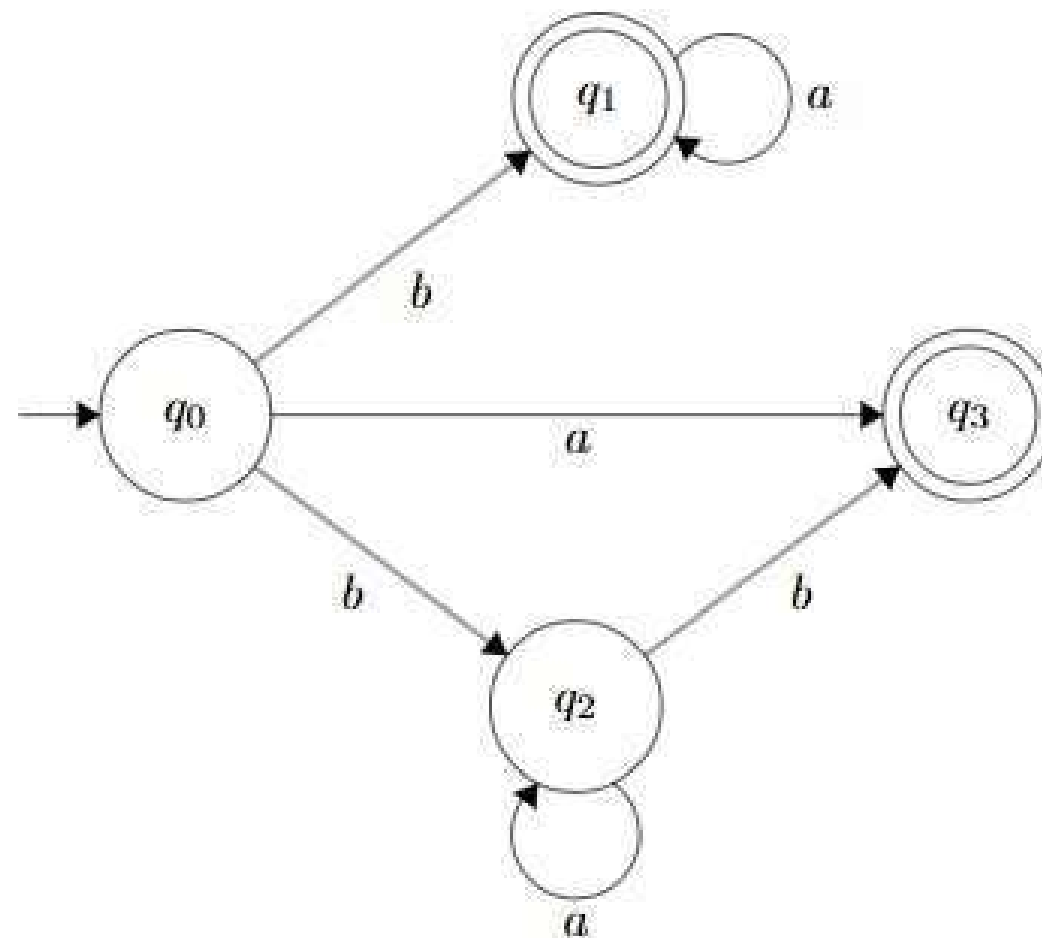
$$\begin{aligned} R &= (aa \cup baa \cup bb), P = (a \cup ab \cup bab) \\ \epsilon &\notin R, \text{ then } \Rightarrow S = R^*P = (aa \cup baa \cup bb)^*(a \cup ab \cup bab) \\ &\Rightarrow \end{aligned}$$

$$(aa|bb|baa)^*(a|ab|bab)$$

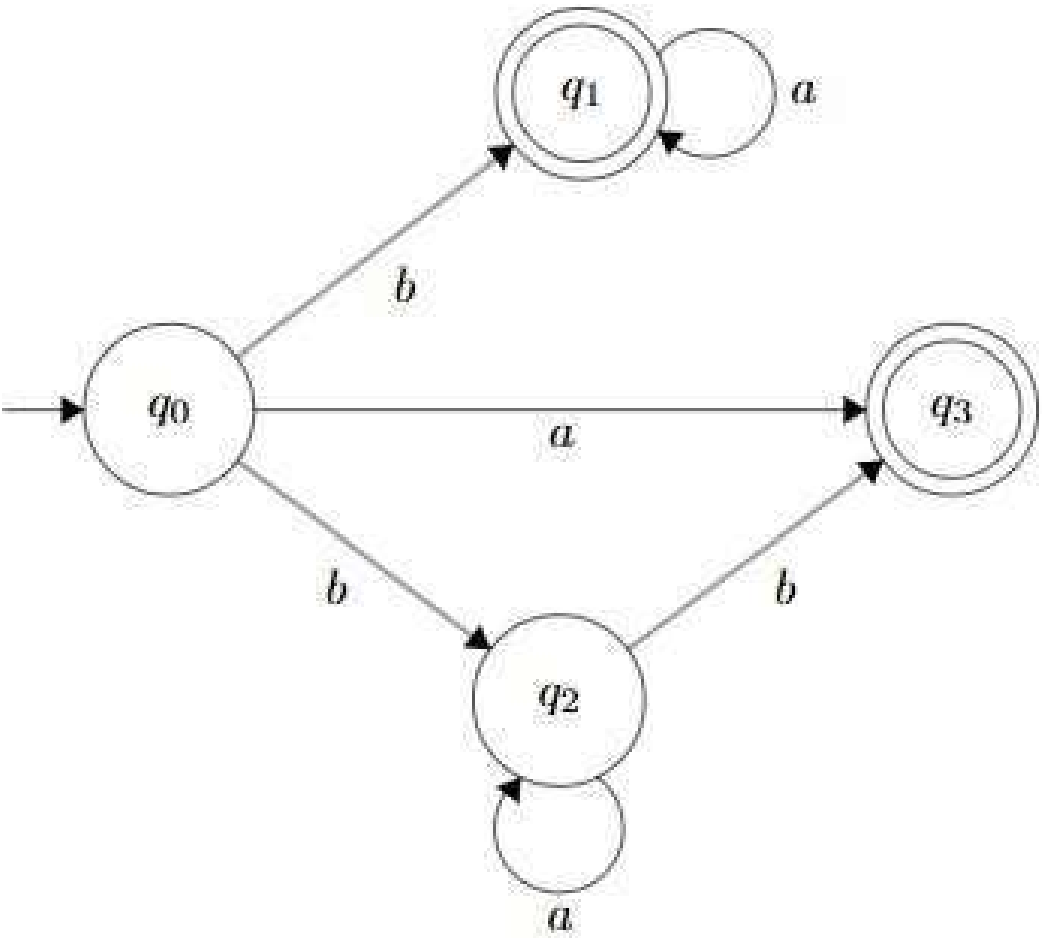
DFA-NFA Equivalency

Question

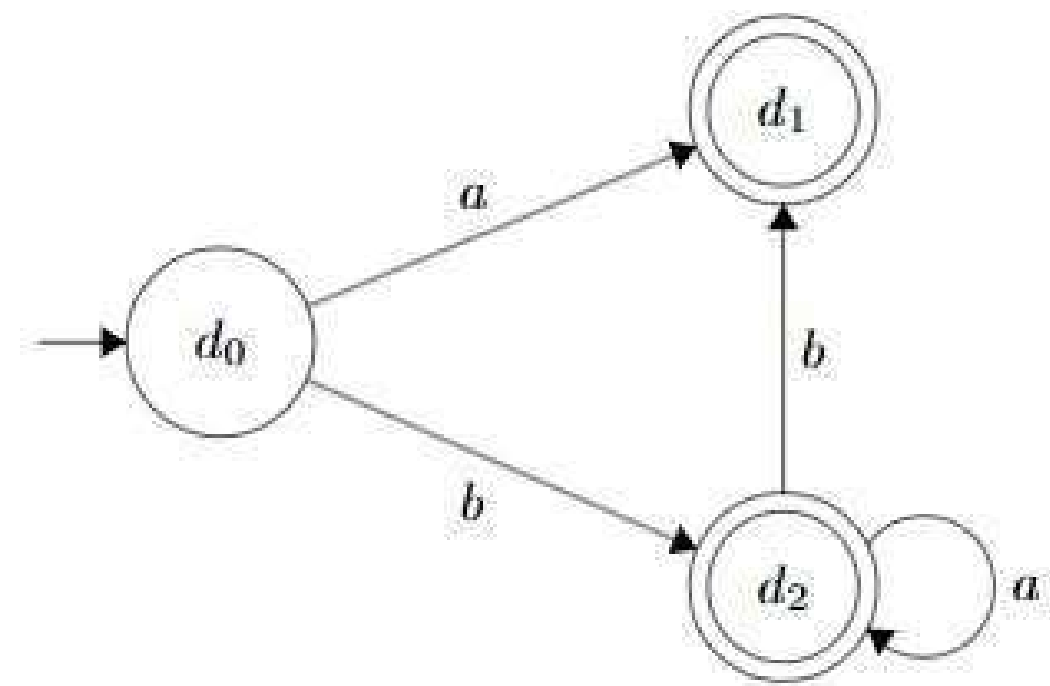
Find the DFA equivalent of the given NFA below.



New State	Next State with 'a'	Next State with 'b'
$d_0 \{q_0\}$	$d_1 \{q_3\}$	$d_2 \{q_1, q_2\}$
$d_1 \{q_3\}$	X	X
$d_2 \{q_1, q_2\}$	$d_2 \{q_1, q_2\}$	$d_1 \{q_3\}$



New State	Next State with 'a'	Next State with 'b'
$d_0 \{q_0\}$	$d_1 \{q_3\}$	$d_2 \{q_1, q_2\}$
$d_1 \{q_3\}$	X	X
$d_2 \{q_1, q_2\}$	$d_2 \{q_1, q_2\}$	$d_1 \{q_3\}$



Question

S $\rightarrow$ aaY | W

X $\rightarrow$ Y | aW

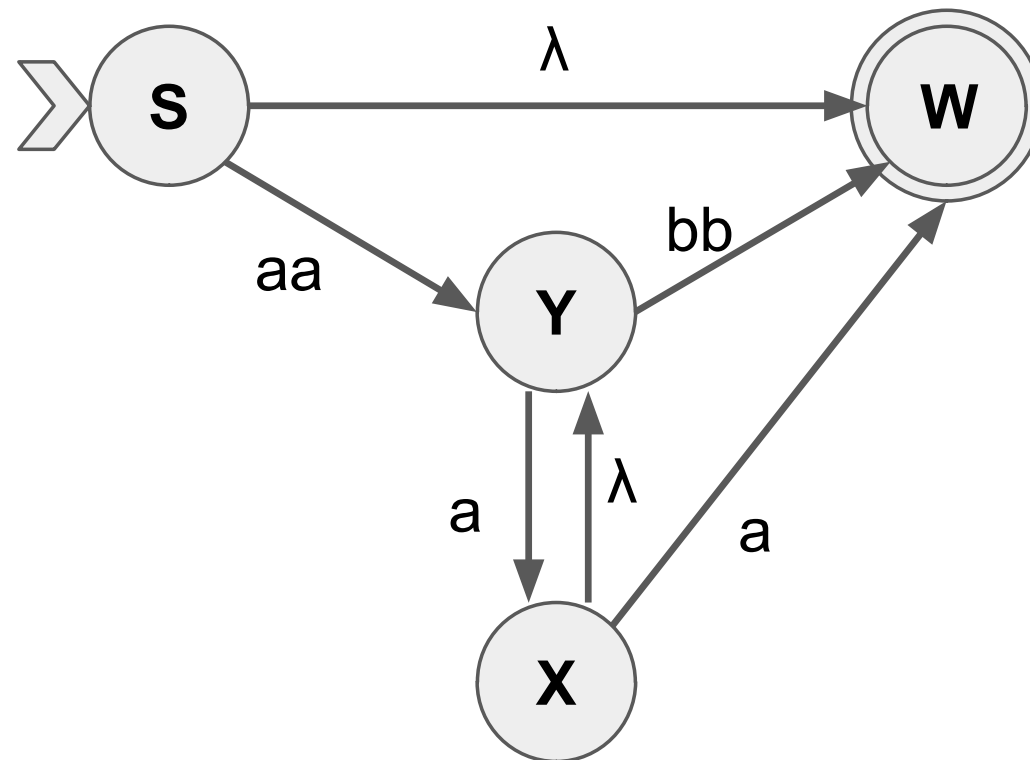
Y $\rightarrow$ aX | bbW

W $\rightarrow$ λ

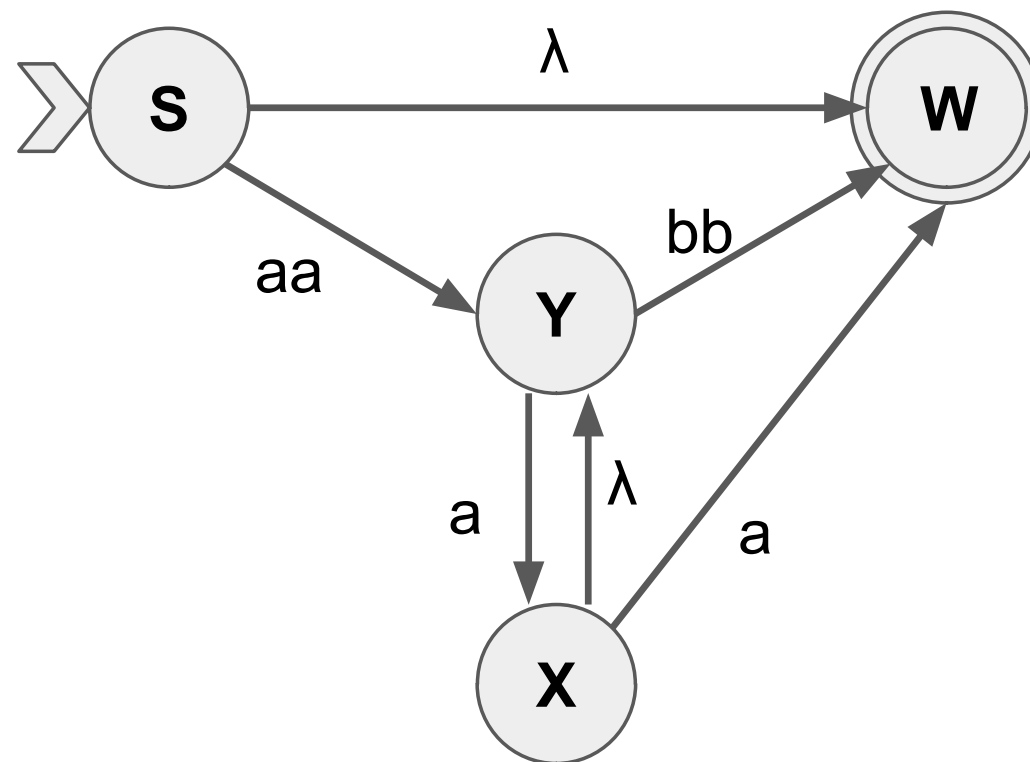
For the given grammar above;

- Draw the NFA that accepts the same language.
- Heuristically, find the regular expression.
- Find the DFA equivalent.

$\mathbf{S} \rightarrow aaY \mid W$
 $\mathbf{X} \rightarrow Y \mid aW$
 $\mathbf{Y} \rightarrow aX \mid bbW$
 $\mathbf{W} \rightarrow \lambda$

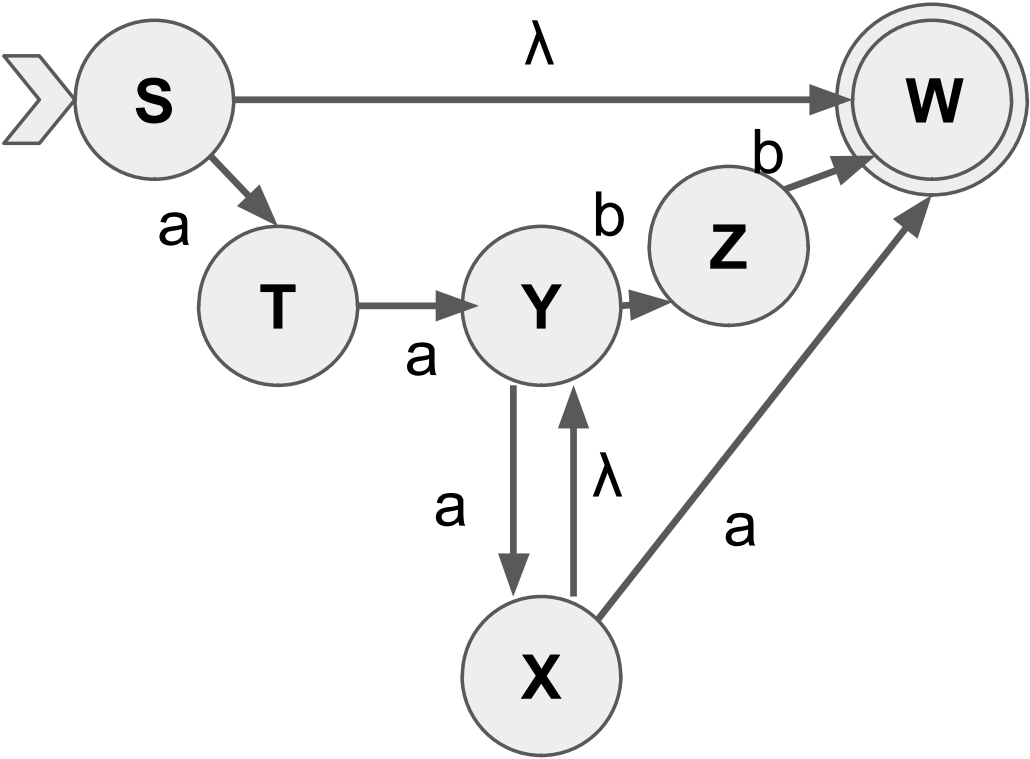


$S \rightarrow aaY \mid W$
 $X \rightarrow Y \mid aW$
 $Y \rightarrow aX \mid bbW$
 $W \rightarrow \lambda$

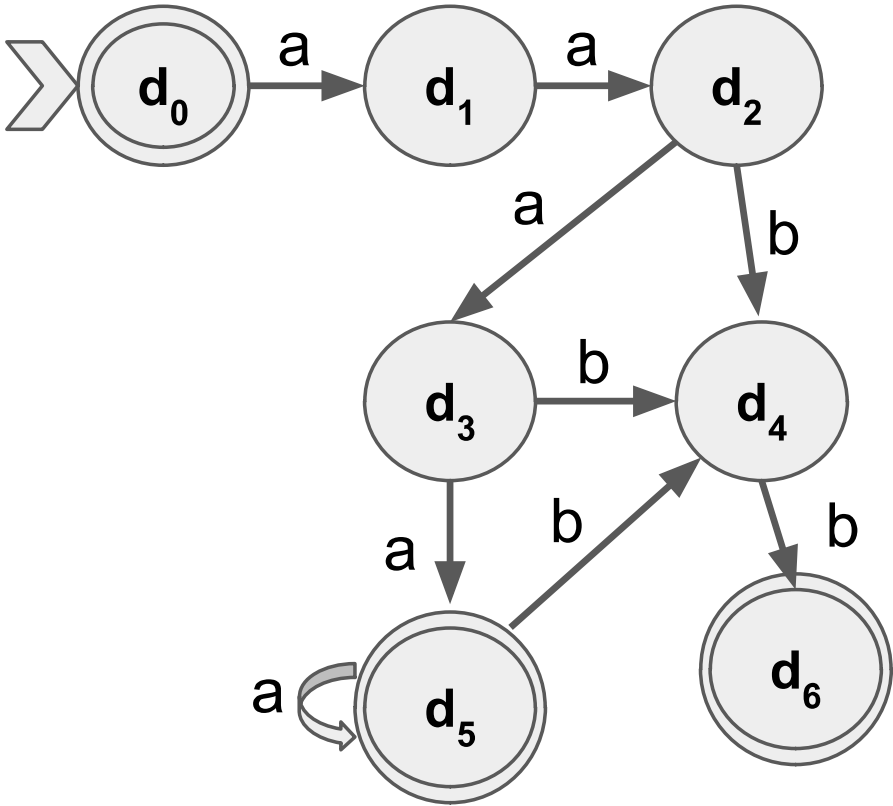


$\lambda \mid aaa^*(aa|bb)$

New State	Next State with 'a'	Next State with 'b'
$d_0 \{S, W\}$	$d_1 \{T\}$	X
$d_1 \{T\}$	$d_2 \{Y\}$	X
$d_2 \{Y\}$	$d_3 \{X, Y\}$	$d_4 \{Z\}$
$d_3 \{X, Y\}$	$d_5 \{X, Y, W\}$	$d_4 \{Z\}$
$d_4 \{Z\}$	X	$d_6 \{W\}$
$d_5 \{X, Y, W\}$	$d_5 \{X, Y, W\}$	$d_4 \{Z\}$
$d_6 \{W\}$	X	X



New State	Next State with 'a'	Next State with 'b'
$d_0 \{S,W\}$	$d_1 \{T\}$	X
$d_1 \{T\}$	$d_2 \{Y\}$	X
$d_2 \{Y\}$	$d_3 \{X,Y\}$	$d_4 \{Z\}$
$d_3 \{X,Y\}$	$d_5 \{X,Y,W\}$	$d_4 \{Z\}$
$d_4 \{Z\}$	X	$d_6 \{W\}$
$d_5 \{X,Y,W\}$	$d_5 \{X,Y,W\}$	$d_4 \{Z\}$
$d_6 \{W\}$	X	X



Question

S $\rightarrow$ 0A | 1B | 0S

A $\rightarrow$ 1B | 0S | 1

B $\rightarrow$ 1B | 0C | 0 | 1

C $\rightarrow$ 1C | 1

For the given grammar above;

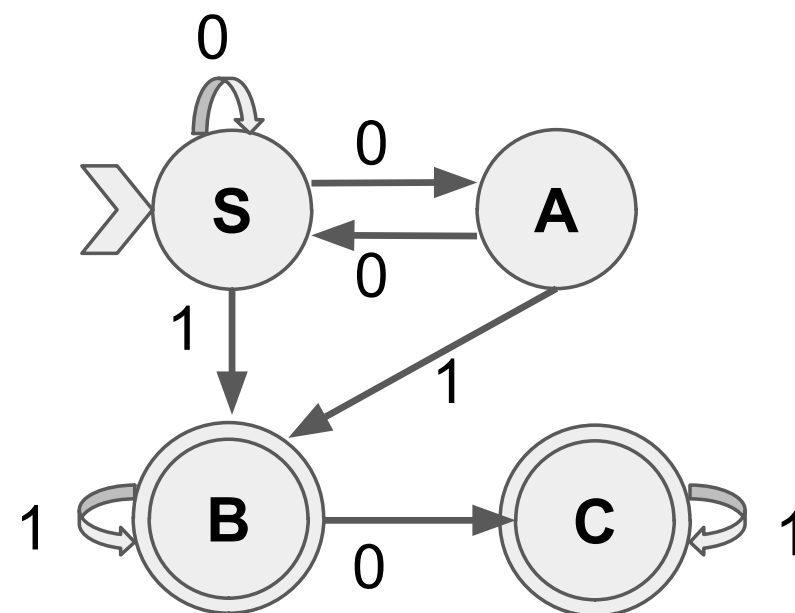
- Draw the NFA that accepts the same language.
- Find the DFA equivalent.

S $\rightarrow$ 0A | 1B | 0S

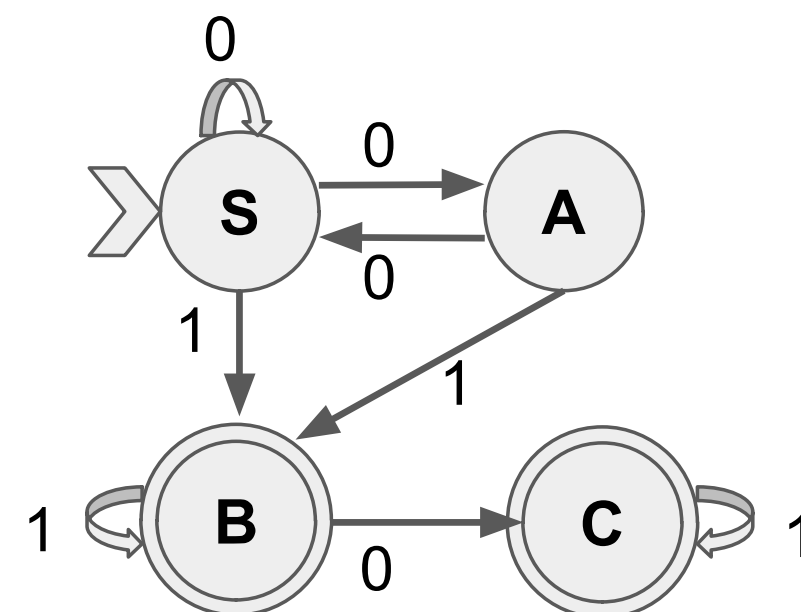
A $\rightarrow$ 1B | 0S | 1

B $\rightarrow$ 1B | 0C | 0 | 1

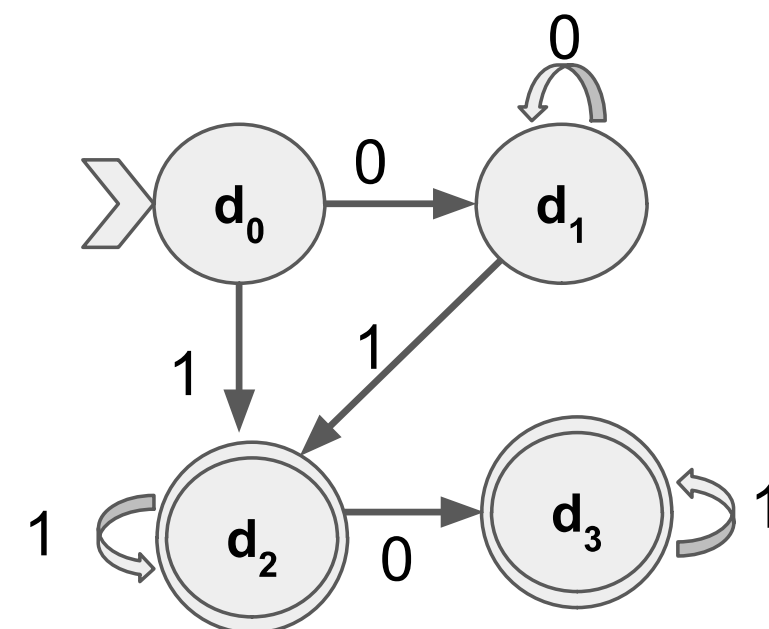
C $\rightarrow$ 1C | 1



New State	Next State with '0'	Next State with '1'
$d_0 \{S\}$	$d_1 \{S,A\}$	$d_2 \{B\}$
$d_1 \{S,A\}$	$d_1 \{S,A\}$	$d_2 \{B\}$
$d_2 \{B\}$	$d_3 \{C\}$	$d_2 \{B\}$
$d_3 \{C\}$	X	$d_3 \{C\}$



New State	Next State with '0'	Next State with '1'
$d_0 \{S\}$	$d_1 \{S,A\}$	$d_2 \{B\}$
$d_1 \{S,A\}$	$d_1 \{S,A\}$	$d_2 \{B\}$
$d_2 \{B\}$	$d_3 \{C\}$	$d_2 \{B\}$
$d_3 \{C\}$	X	$d_3 \{C\}$

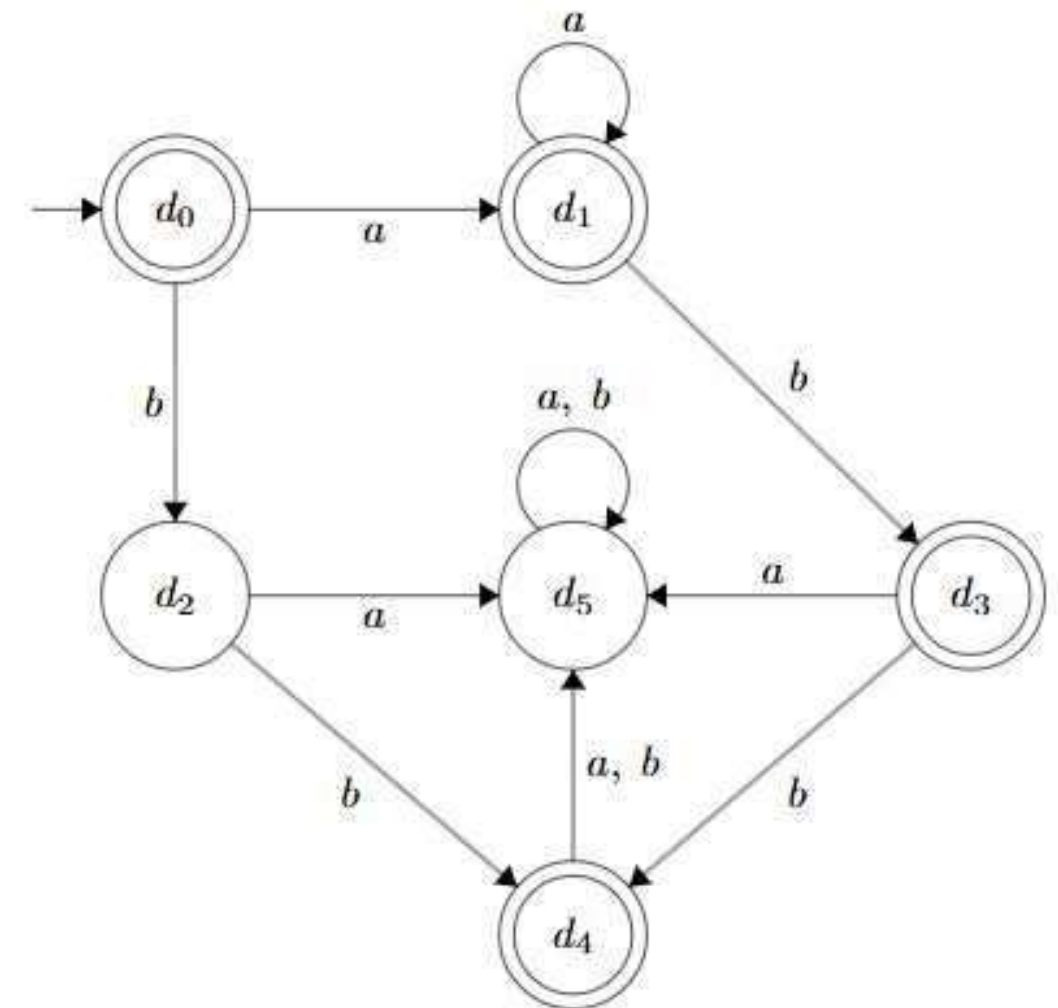


Recognizing Regular Expressions

Question

For the given automata;

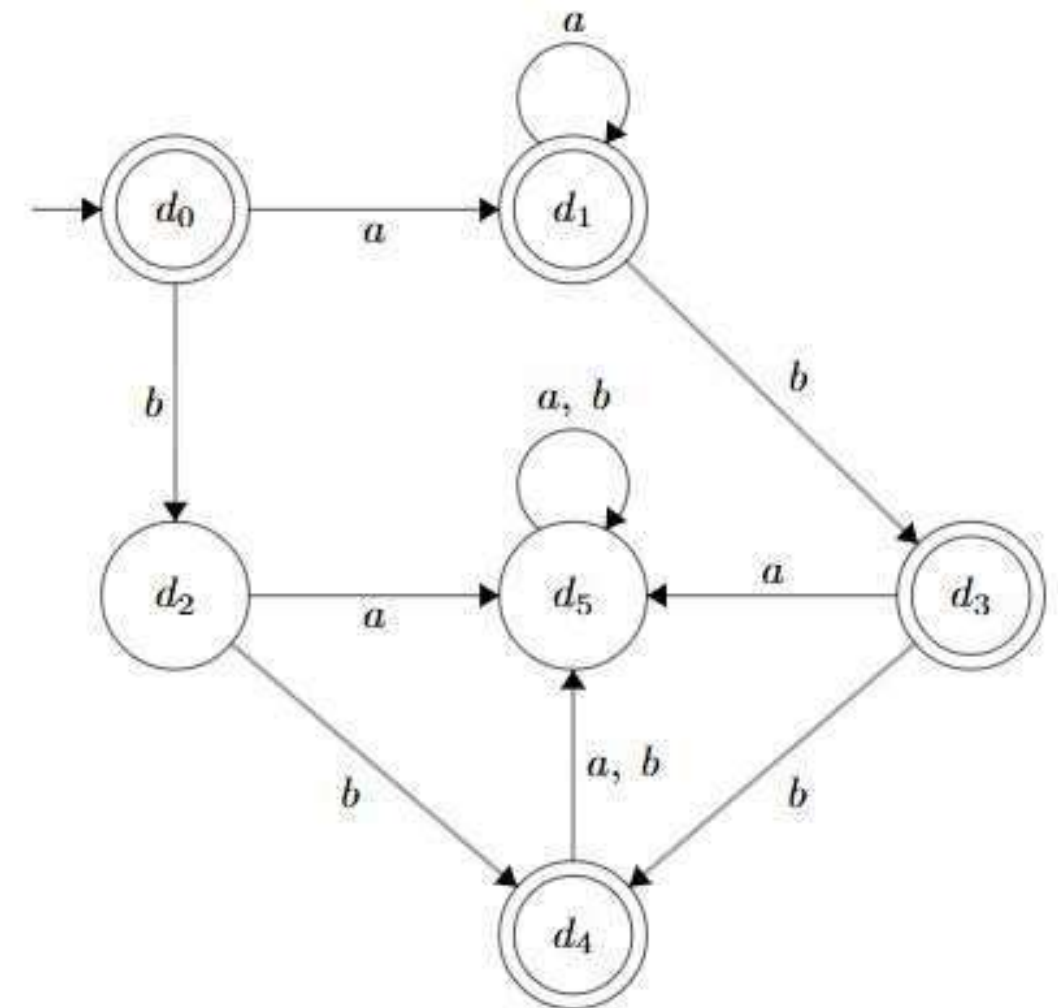
- Heuristically, find it's regular expression.
- Systematically create the regular expression and compare it to the previously defined expression.



No loop, we can check paths to accepting states

- For $d_0 \rightarrow \lambda$
- For $d_1 \rightarrow a^+$
- For $d_3 \rightarrow a^+b$
- For $d_4 \rightarrow bb$ or a^+bb

Thus we can state: $\text{Regex} \rightarrow a^* \vee a^+b \vee a^+bb$

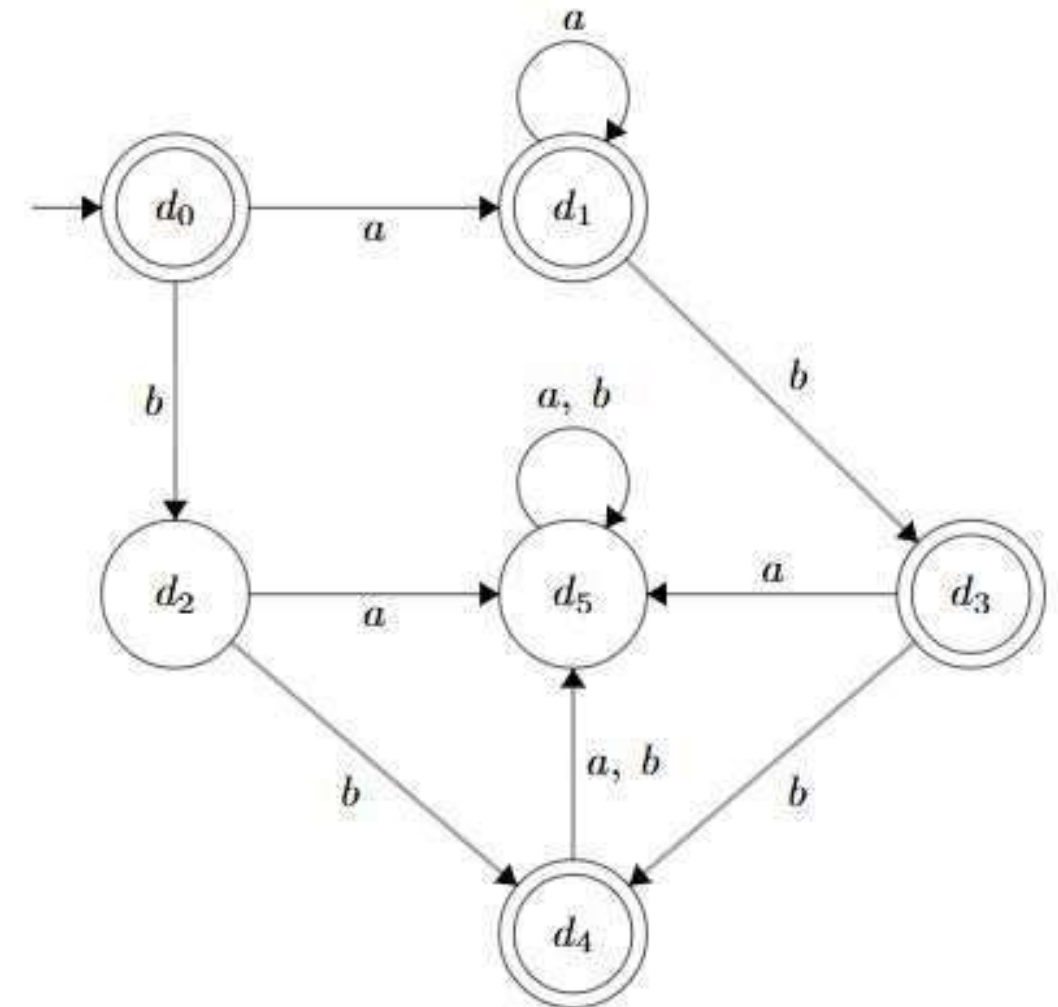


- $d_0 = \lambda$
- $d_1 = d_0 a \vee d_1 a$
- $d_2 = d_0 b$
- $d_3 = d_1 b$
- $d_4 = d_2 b \vee d_3 b$

ReGex will be $d_0 \vee d_1 \vee d_3 \vee d_4$

- $d_0 \rightarrow \lambda = \lambda$
- $d_1 \rightarrow a \vee d_1 a = aa^*$
- $d_2 \rightarrow d_0 b = b$
- $d_3 \rightarrow d_1 b = aa^*b$
- $d_4 \rightarrow d_2 b \vee d_3 b = bb \vee aa^*bb$

$$d_0 \vee d_1 \vee d_3 \vee d_4 = \lambda \vee a^+ \vee a^+b \vee bb \vee a^+bb = a^* \vee a^+b \vee a^*bb$$

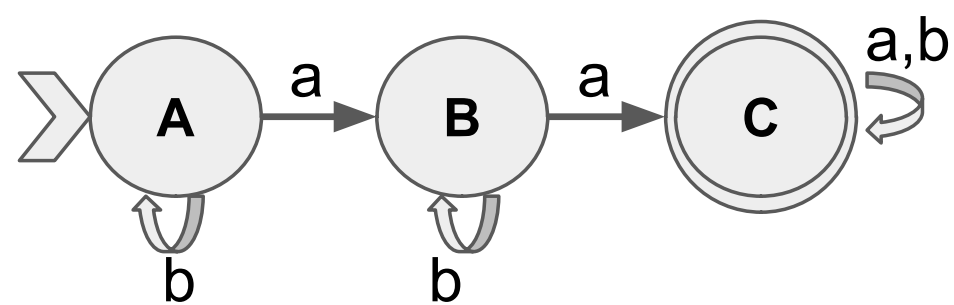


Question

$L = \{w: w \text{ has at least two a's } \mathbf{OR} \text{ an odd number of b's}\}$ is given.

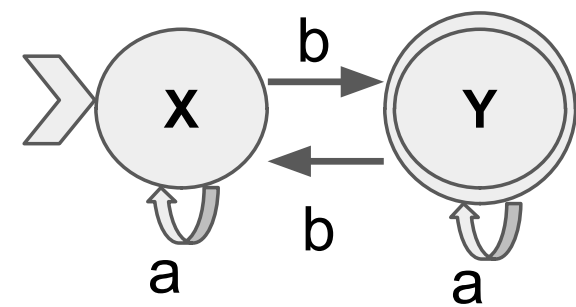
- Design a DFA recognizing both parts of the given language separately.
- Systematically combine two DFAs.

L1= {w: w has at least two a's}



Current State	Next State with 'a'	Next State with 'b'
A	B	A
B	C	B
C	C	C

L2= {w: w has an odd number of b's}



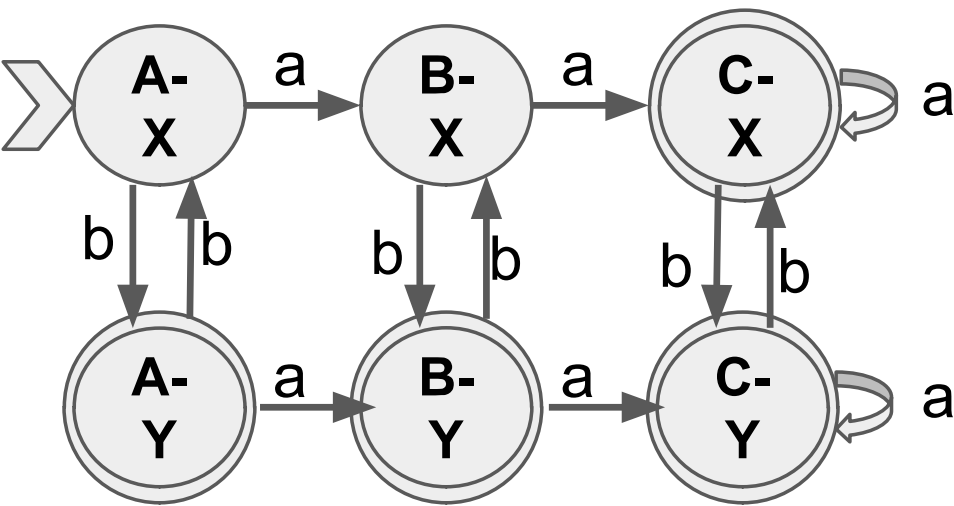
Current State	Next State with 'a'	Next State with 'b'
X	X	Y
Y	Y	X

Current State	Next State with 'a'	Next State with 'b'
A	B	A
B	C	B
C	C	C

Current State	Next State with 'a'	Next State with 'b'
X	X	Y
Y	Y	X

New State	Next State with 'a'	Next State with 'b'
A-X	B-X	A-Y
B-X	C-X	B-Y
C-X	C-X	C-Y
A-Y	B-Y	A-X
B-Y	C-Y	B-X
C-Y	C-Y	C-X

New State	Next State with 'a'	Next State with 'b'
A-X	B-X	A-Y
B-X	C-X	B-Y
C-X	C-X	C-Y
A-Y	B-Y	A-X
B-Y	C-Y	B-X
C-Y	C-Y	C-X



$$\begin{aligned} &ba(a+aa^*(a+\$+b(a+ba^*b)^*(b+ba^*(a+\$)))) \\ &+(a+bab)(bb)^*(b(a+aa^*(a+\$))+(a+baa^*b)(\\ &a+ba^*b)^*(b+ba^*(a+\$)))+(\$+bb)(bb)^*(ba(a \\ &+aa^*(a+\$+b(a+ba^*b)^*(b+ba^*(a+\$))))+(a+ \\ &bab)(bb)^*(b(a+aa^*(a+\$))+(a+baa^*b)(a+b \\ &a^*b)^*(b+ba^*(a+\$)))) \end{aligned}$$

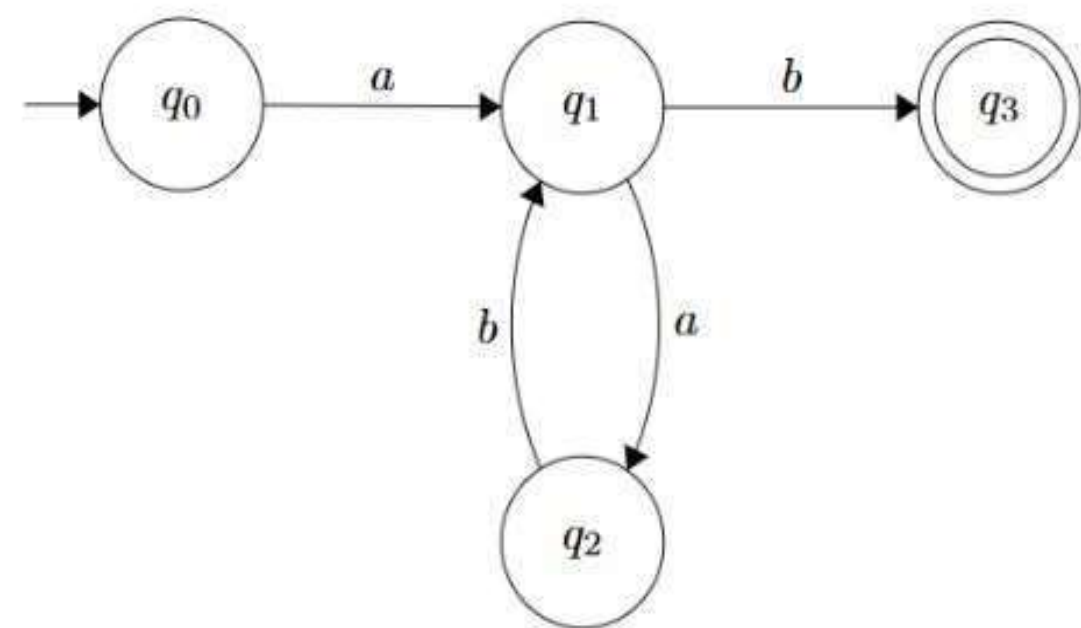
Question

Using a systematic approach, check whether each of the following regular expression couples are disjoint, and if they are not, find the shortest non-empty (not λ) string that matches both of them

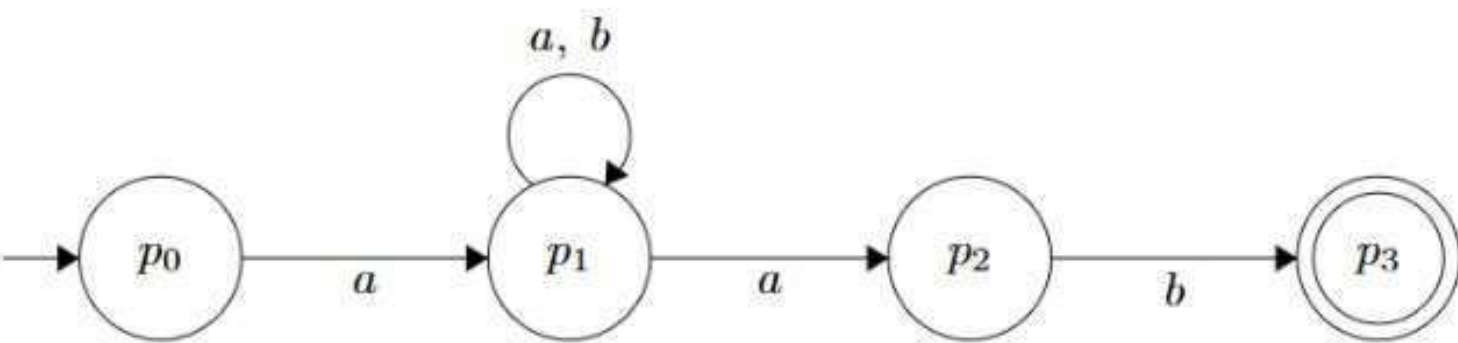
A. $L1 = a(ab)^*b$ and $L2 = a(a \vee b)^*ab$

B. $L1 = a(ab)^*a$ and $L2 = a(a \vee b)^*ba$

L1 = a(ab)*b and L2 = a(a v b)*ab



L1



L2

Current State	Next State with 'a'	Next State with 'b'
q0	q1	X
q1	q2	q3
q2	X	q1
q3	X	X

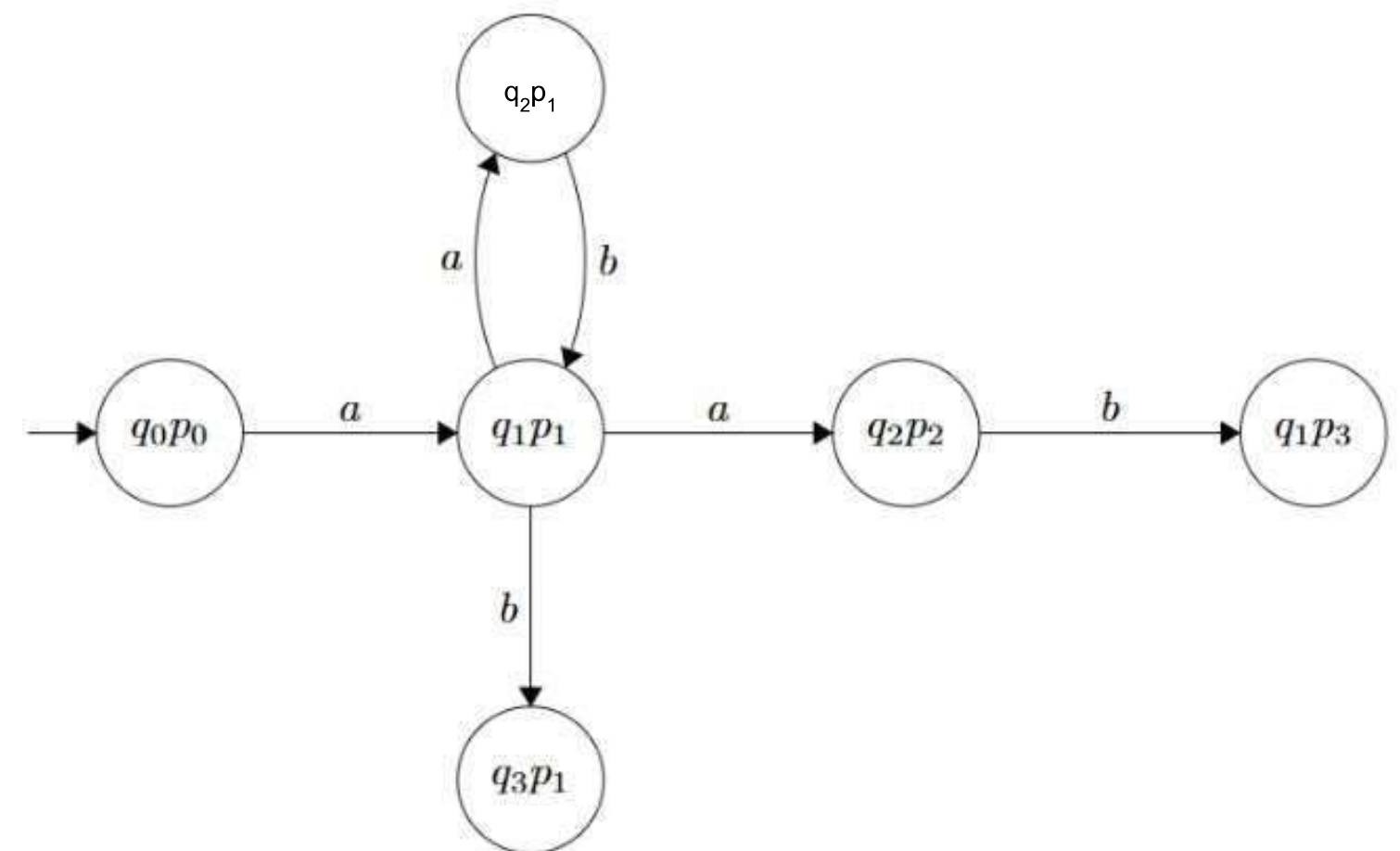
Current State	Next State with 'a'	Next State with 'b'
p0	p1	X
p1	p1,p2	p1
p2	X	p3
p3	X	X

Current State	Next State with 'a'	Next State with 'b'
q0	q1	X
q1	q2	q3
q2	X	q1
q3	X	X

Current State	Next State with 'a'	Next State with 'b'
p0	p1	X
p1	p1,p2	p1
p2	X	p3
p3	X	X

Current State	Next State with 'a'	Next State with 'b'
q0-p0	q1-p1	X
q1-p1	q2-p1,q2-p2	q3-p1
q2-p1	X	q1-p1
q2-p2	X	q1-p3
q3-p1	X	X
q1-p3	X	X

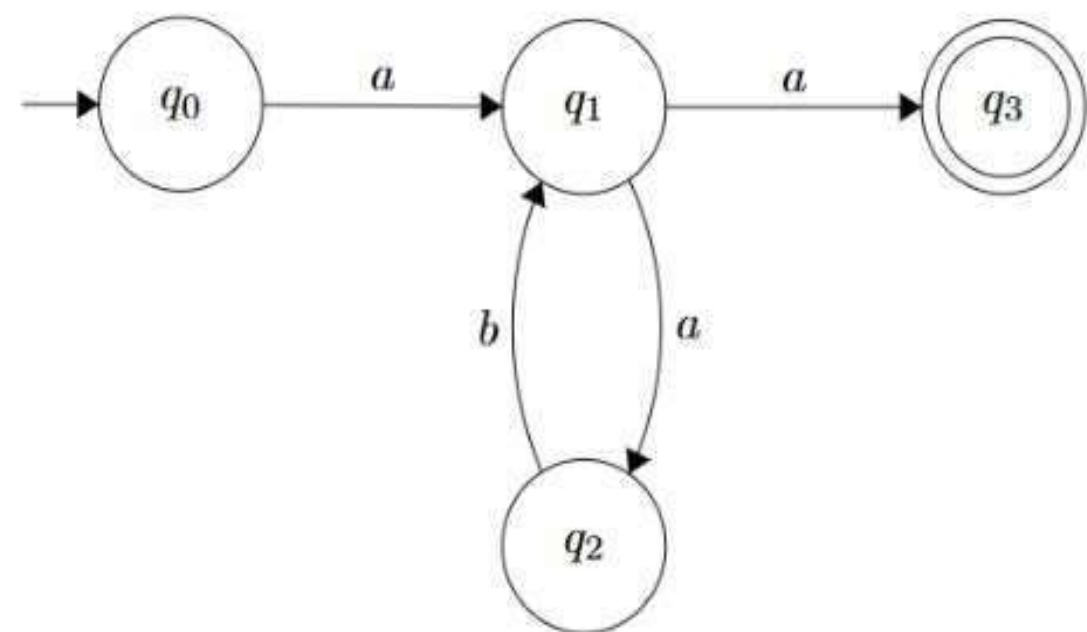
Current State	Next State with 'a'	Next State with 'b'
q0-p0	q1-p1	X
q1-p1	q2-p1,q2-p2	q3-p1
q2-p1	X	q1-p1
q2-p2	X	q1-p3
q3-p1	X	X
q1-p3	X	X



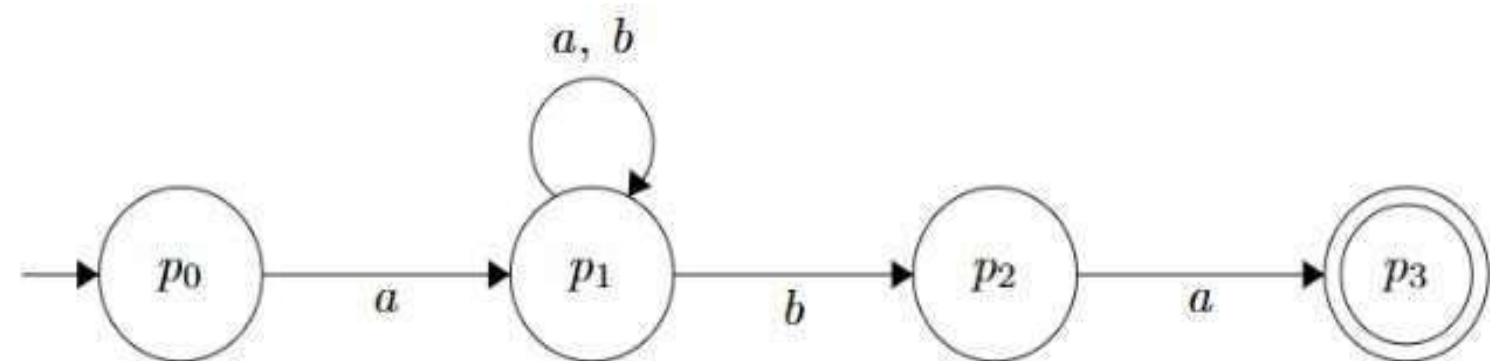
L1 and L2

As there is no reachable final state, we can not find a string

L1 = a(ab)*a and L2 = a(a v b)*ba



L1



L2

Current State	Next State with 'a'	Next State with 'b'
q0	q1	X
q1	q2,q3	X
q2	X	q1
q3	X	X

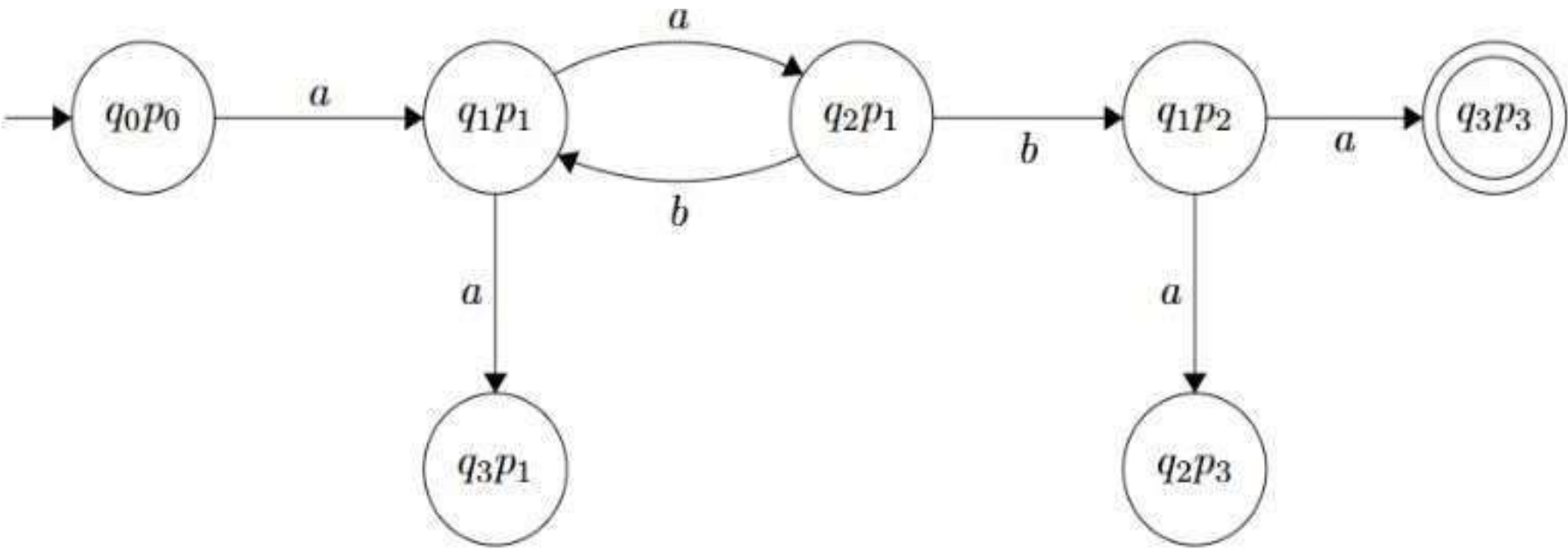
Current State	Next State with 'a'	Next State with 'b'
p0	p1	X
p1	p1	p1,p2
p2	p3	X
p3	X	X

Current State	Next State with 'a'	Next State with 'b'
q0	q1	X
q1	q2,q3	X
q2	X	q1
q3	X	X

Current State	Next State with 'a'	Next State with 'b'
p0	p1	X
p1	p1	p1,p2
p2	p3	X
p3	X	X

Current State	Next State with 'a'	Next State with 'b'
q0-p0	q1-p1	X
q1-p1	q2-p1,q3-p1	X
q2-p1	X	q1-p1,q1-p2
q3-p1	X	X
q1-p2	q2-p3,q3-p3	X
q2-p3	X	X
q3-p3	X	X

Current State	Next State with 'a'	Next State with 'b'
q0-p0	q1-p1	X
q1-p1	q2-p1,q3-p1	X
q2-p1	X	q1-p1,q1-p2
q3-p1	X	X
q1-p2	q2-p3,q3-p3	X
q2-p3	X	X
q3-p3	X	X



L1 and L2

Shortest string: aaba